

# PETRI NETS

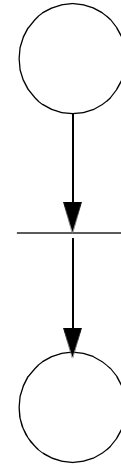
**1. Basic Petri Net Model**

**2. Properties and Analysis of Petri Nets**

**3. Extended Petri Net Models**

# Petri Nets

- **Systems are specified as a directed bipartite graph.**  
**The two kinds of nodes in the graph:**

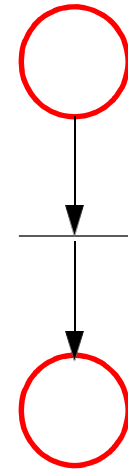


# Petri Nets

- Systems are specified as a directed bipartite graph.

The two kinds of nodes in the graph:

1. Places: they hold the distributed state of the system expressed by the presence/absence of tokens in the places.

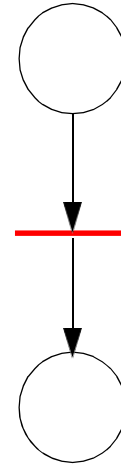


# Petri Nets

- Systems are specified as a directed bipartite graph.

The two kinds of nodes in the graph:

1. Places: they hold the distributed state of the system expressed by the presence/absence of tokens in the places.
2. Transitions: denote the activity in the system

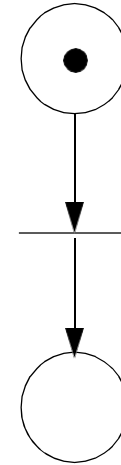


# Petri Nets

- Systems are specified as a directed bipartite graph.

The two kinds of nodes in the graph:

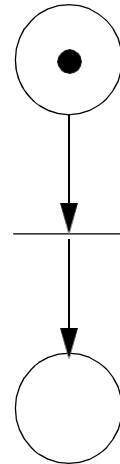
1. Places: they hold the distributed state of the system expressed by the presence/absence of tokens in the places.
2. Transitions: denote the activity in the system



- *The state of the system*: captured by the marking of the places (number of tokens in each place)

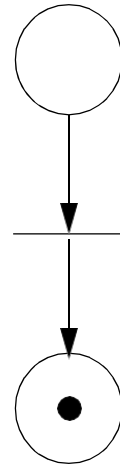
# Petri Nets

- The dynamic evolution of the system: determined by the firing process of transitions.
  - A transition is enabled and may fire whenever all its predecessor places are marked.



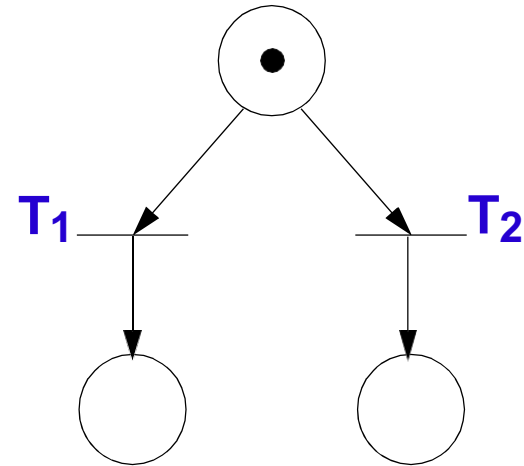
# Petri Nets

- The dynamic evolution of the system: determined by the firing process of transitions.
  - A transition is enabled and may fire whenever all its predecessor places are marked.
  - If a transition fires it removes a token from each predecessor place and adds a token to each successor place.



# Nondeterminism

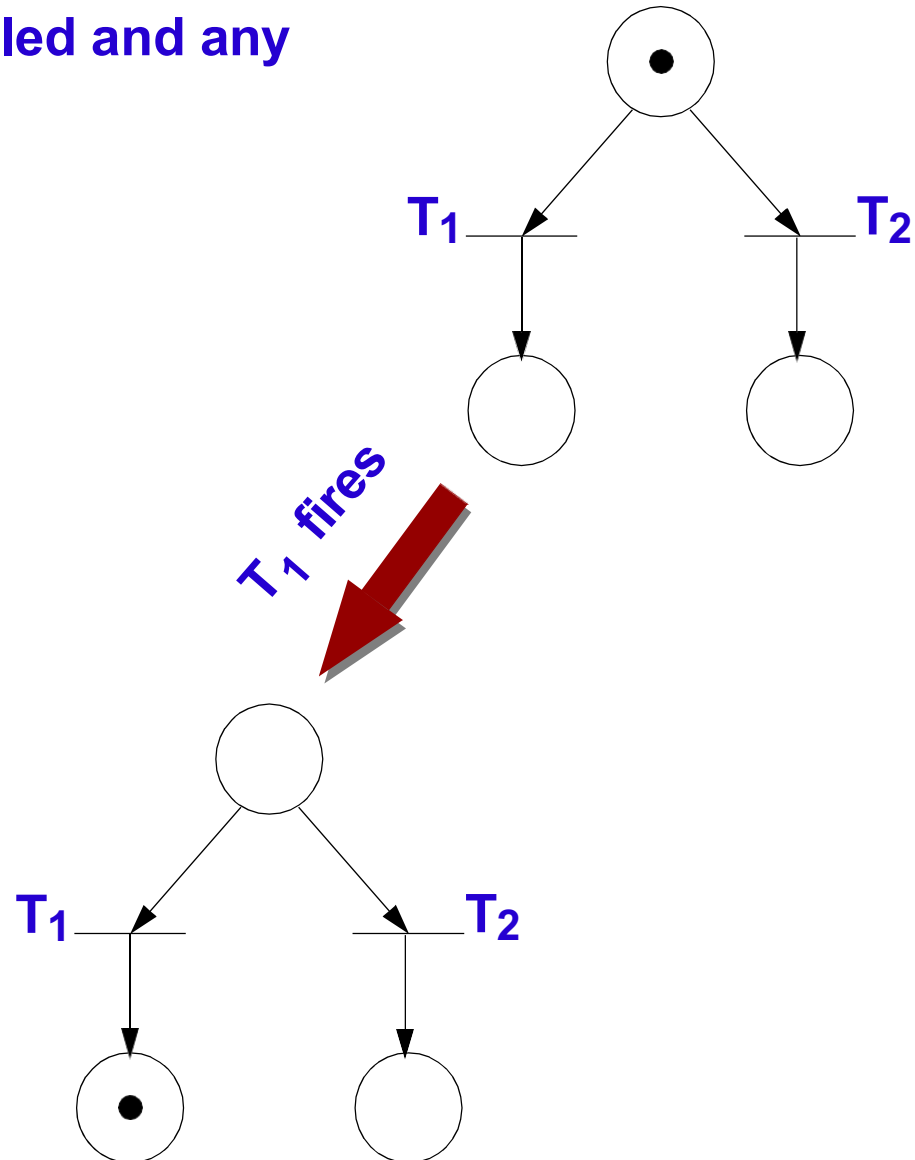
- Both T1 and T2 are enabled and any of the two may fire.





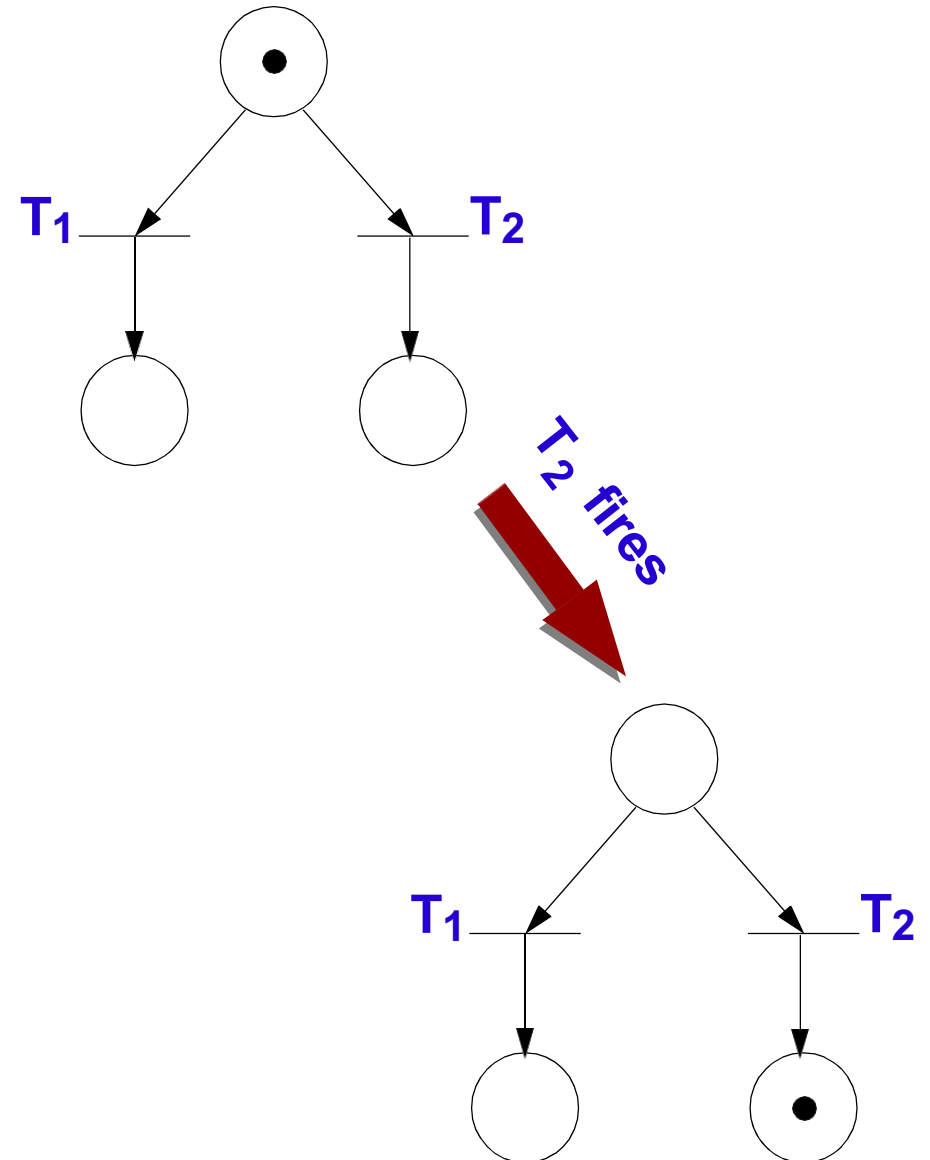
# Nondeterminism

- Both T1 and T2 are enabled and any of the two may fire.



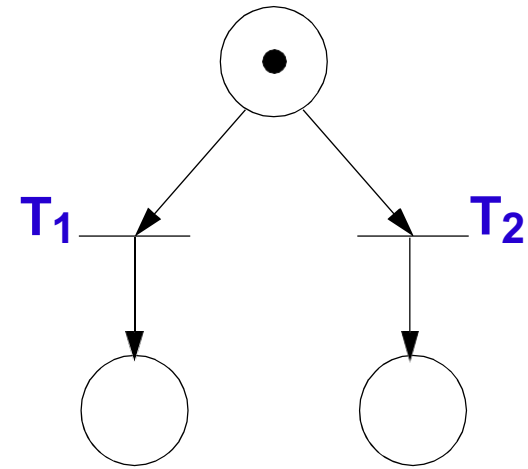
# Nondeterminism

- Both T1 and T2 are enabled and any of the two may fire.

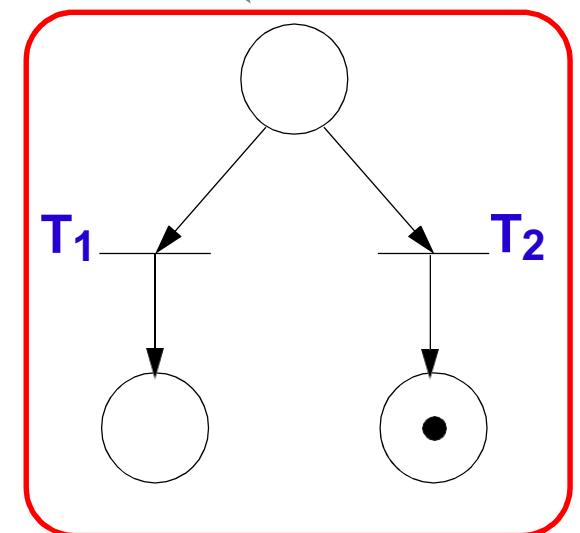
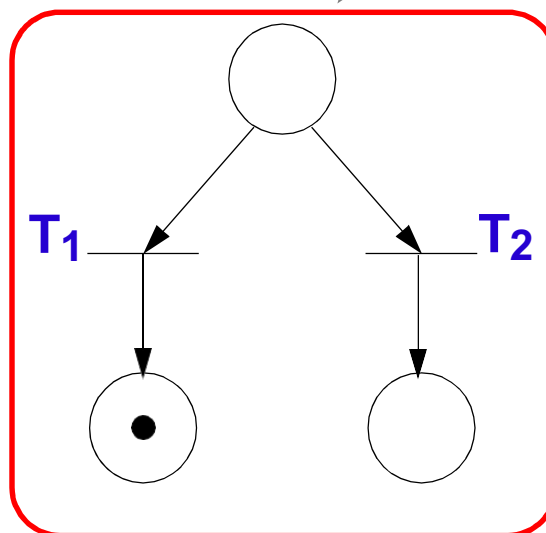


# Nondeterminism

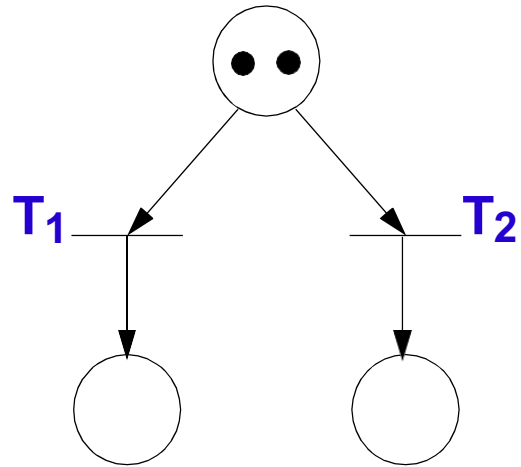
- Both T1 and T2 are enabled and any of the two may fire.



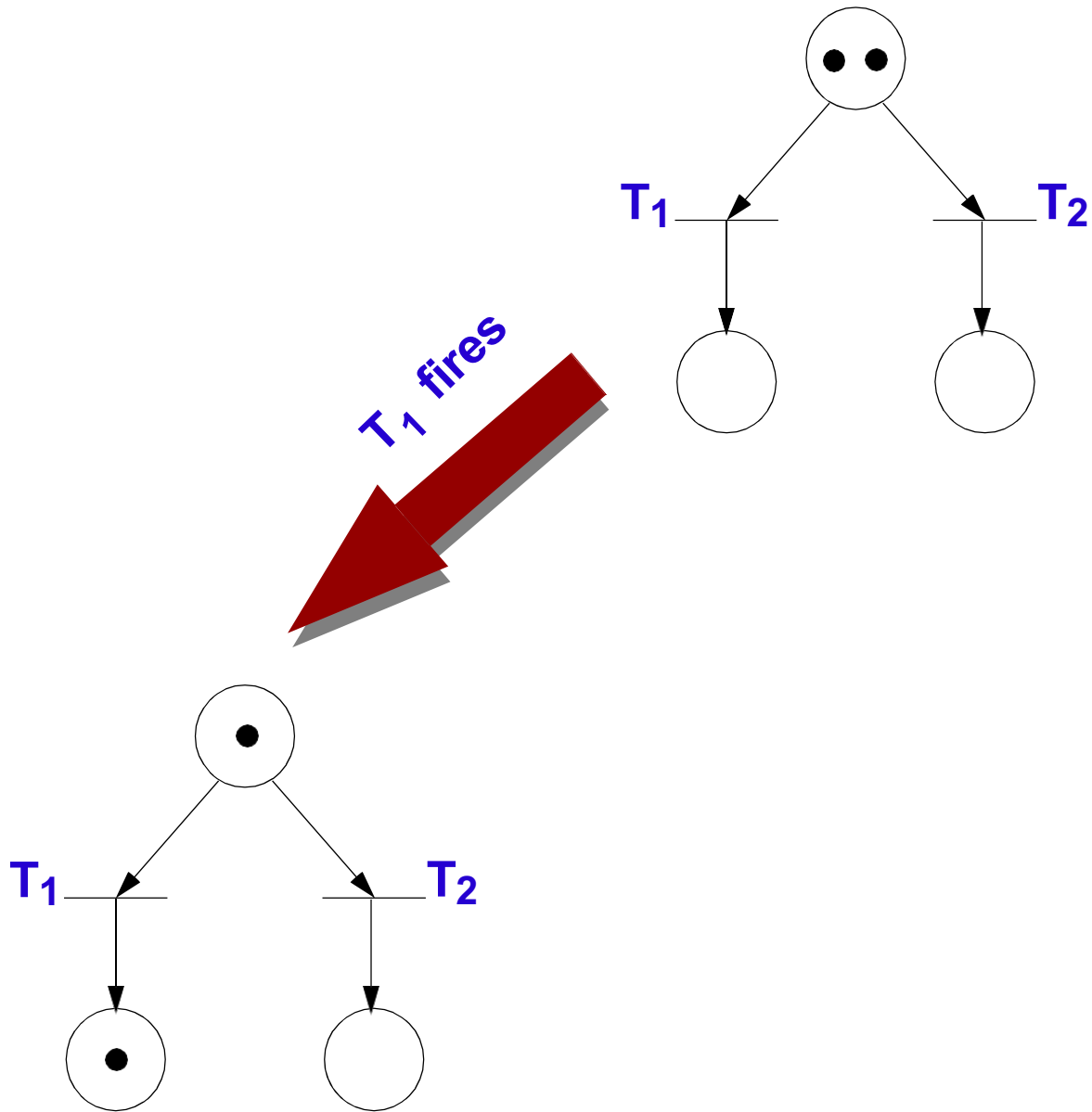
- Any of these two states might be the next state.



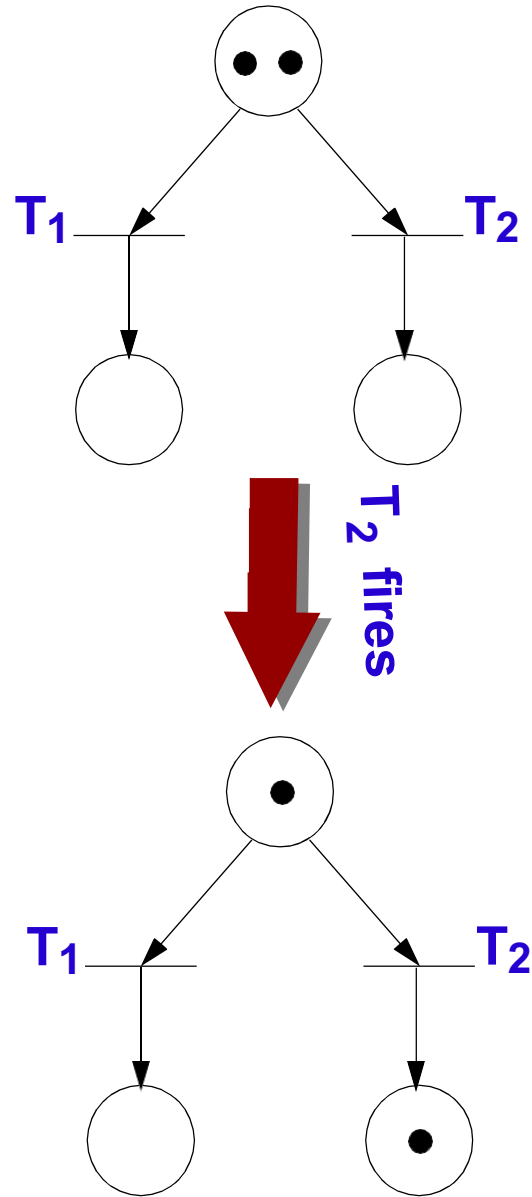
# Nondeterminism



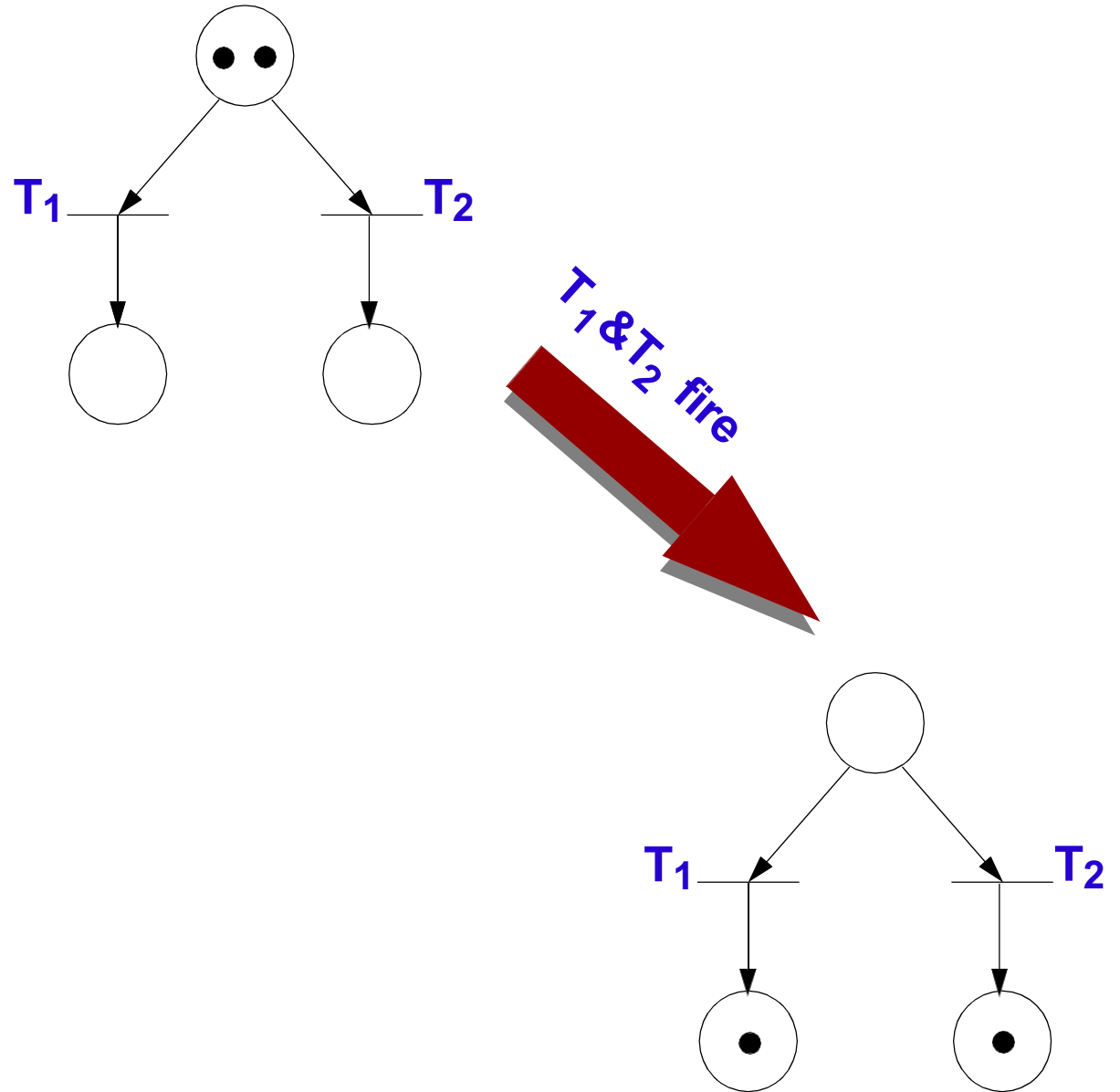
# Nondeterminism



# Nondeterminism

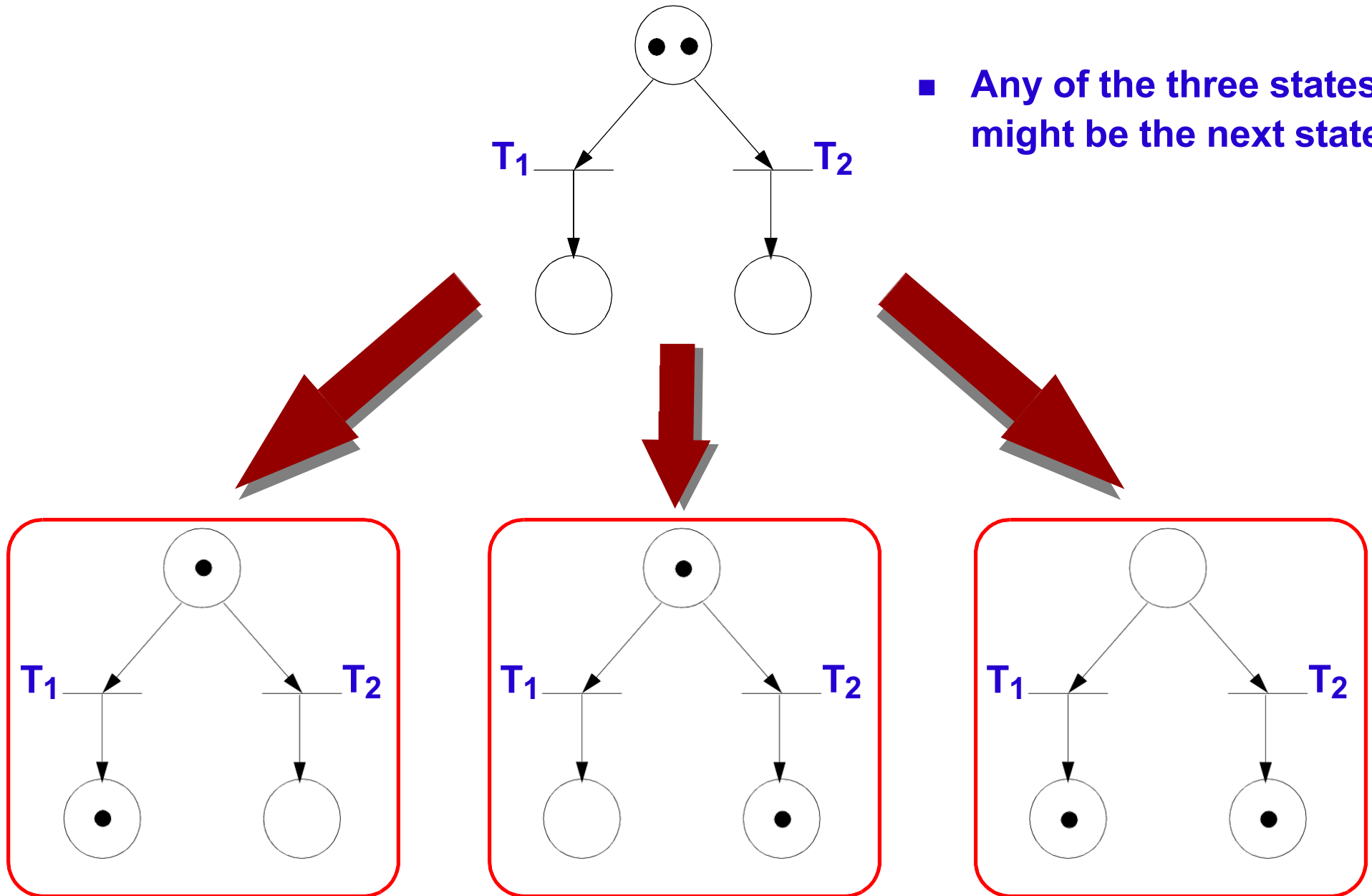


# Nondeterminism



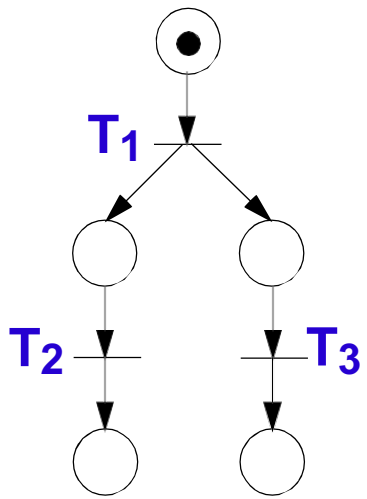
# Nondeterminism

- Any of the three states might be the next state.



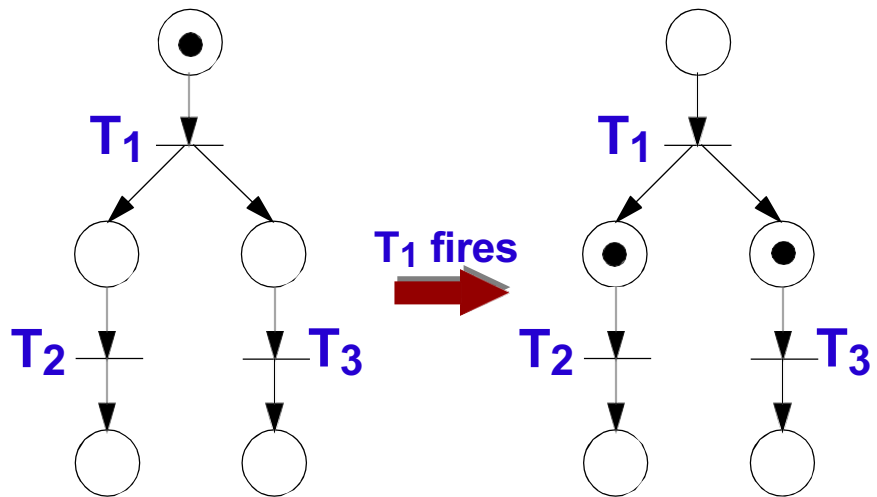


# Nondeterminism

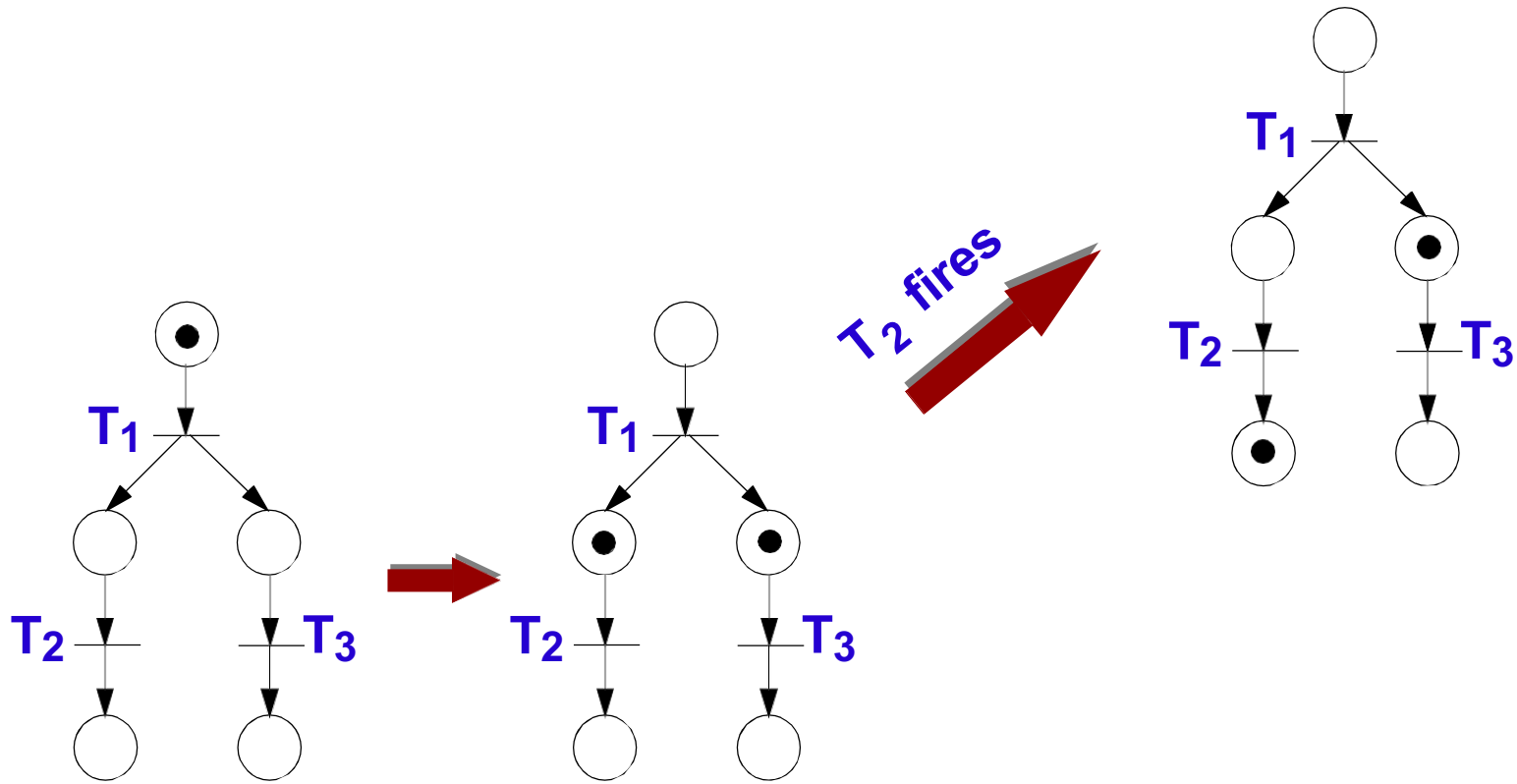


# Nondeterminism

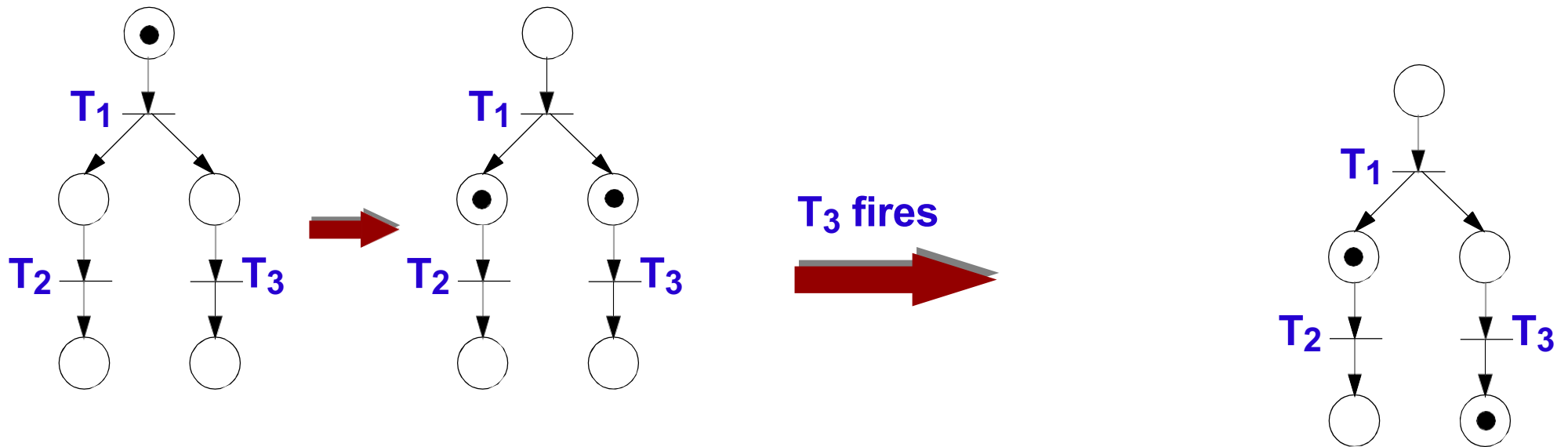
- No nondeterminism here:  
 $T_1$  is the only enabled transition!



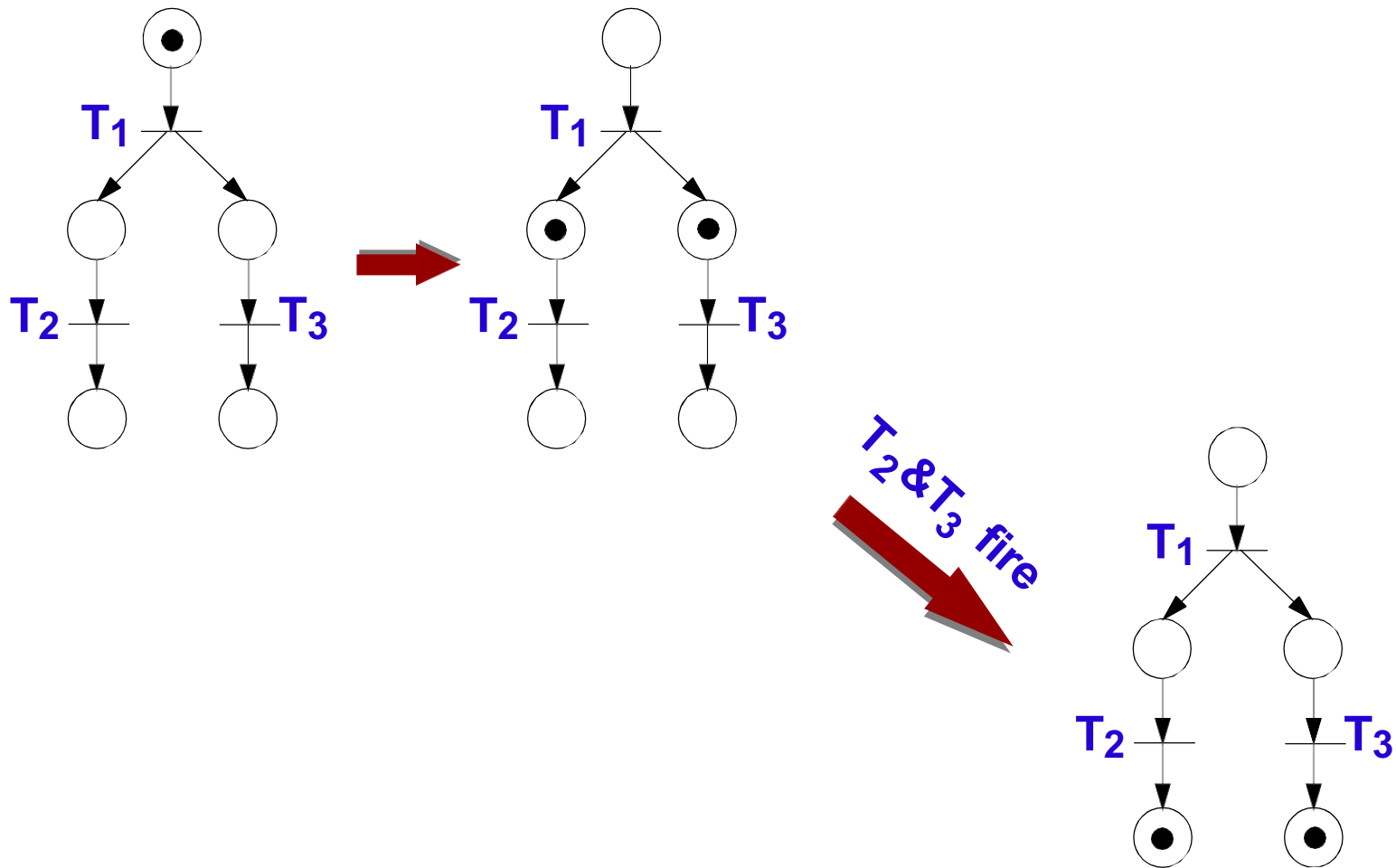
# Nondeterminism



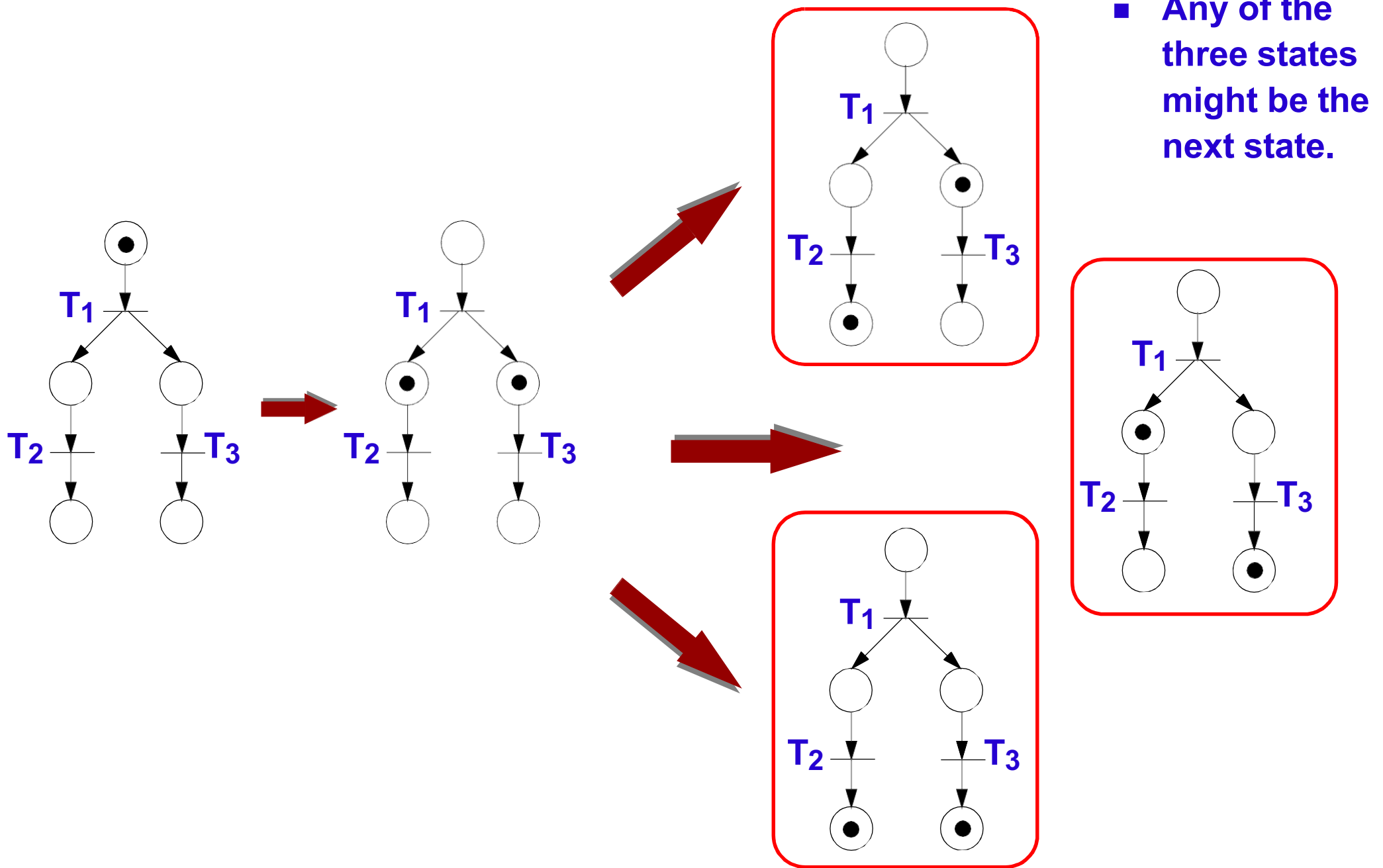
# Nondeterminism



# Nondeterminism



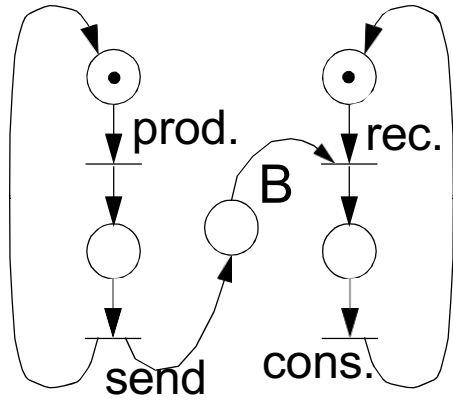
# Nondeterminism



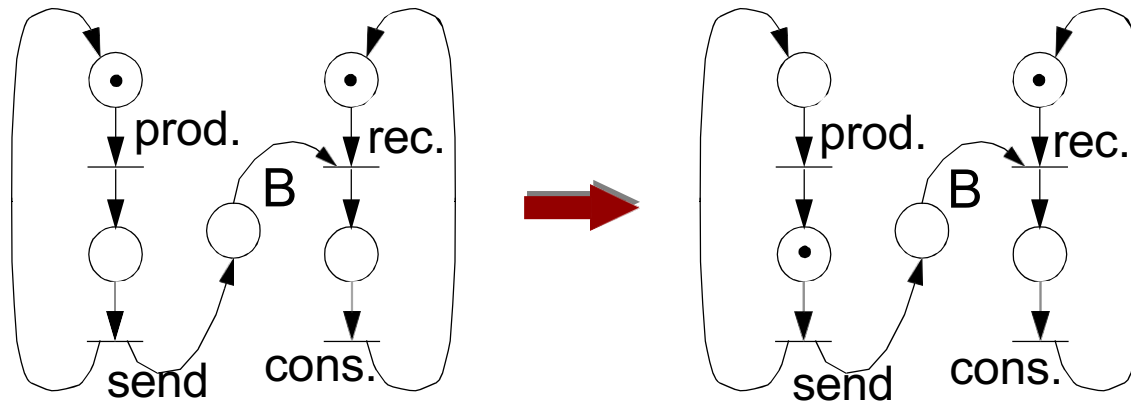
- Any of the three states might be the next state.

# Petri Net Example

A producer and a consumer process communicating through a buffer:

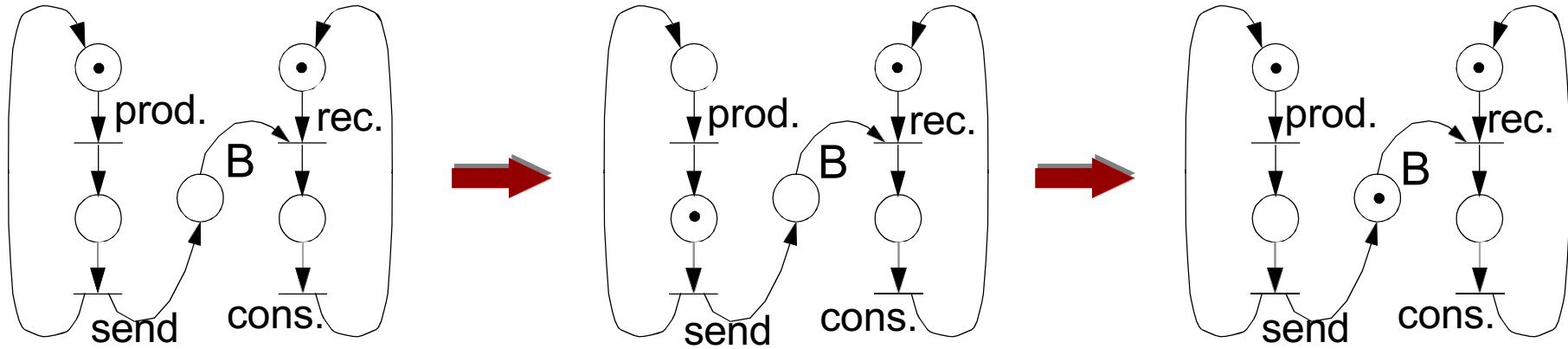


# Petri Net Example

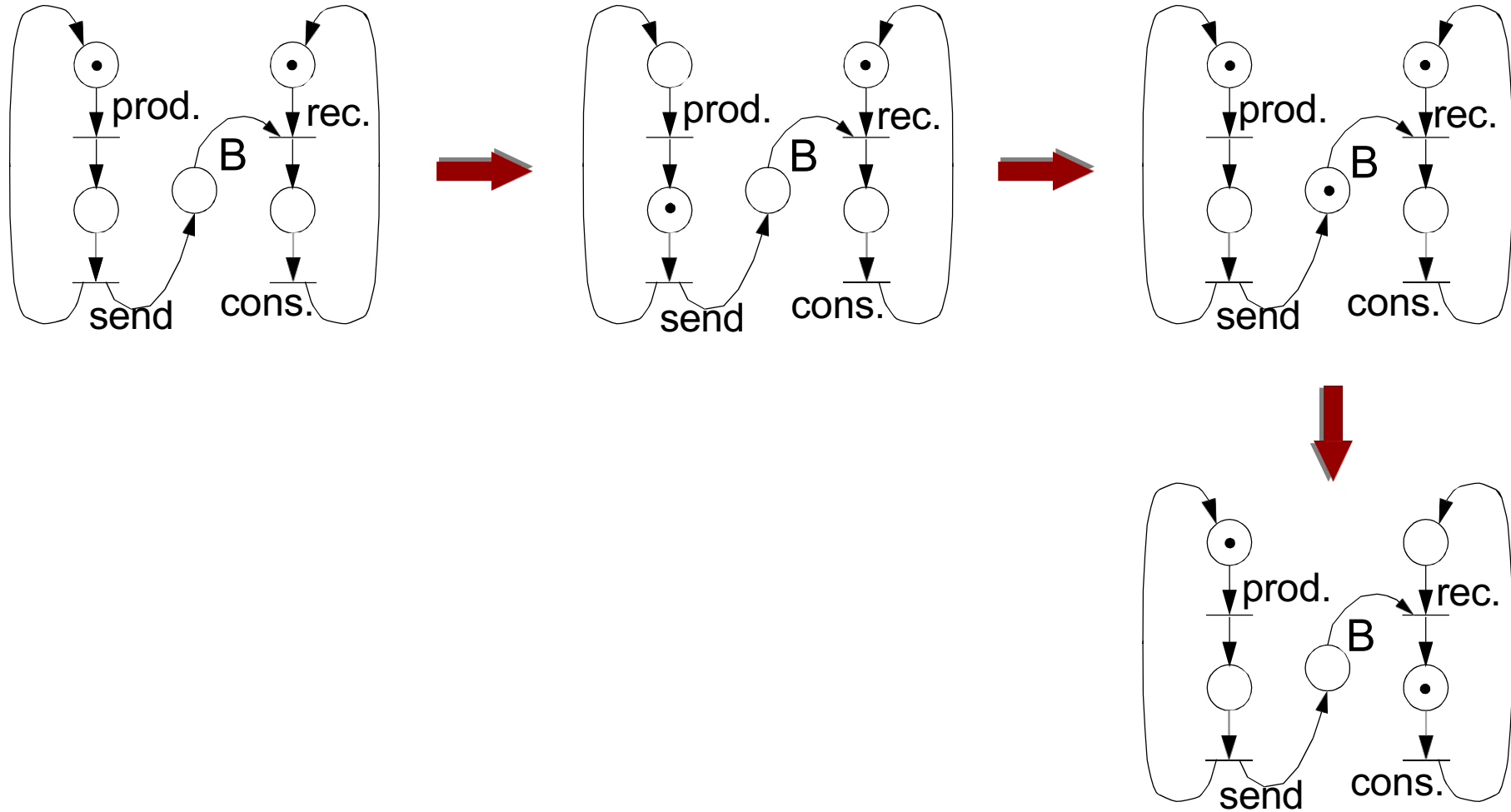




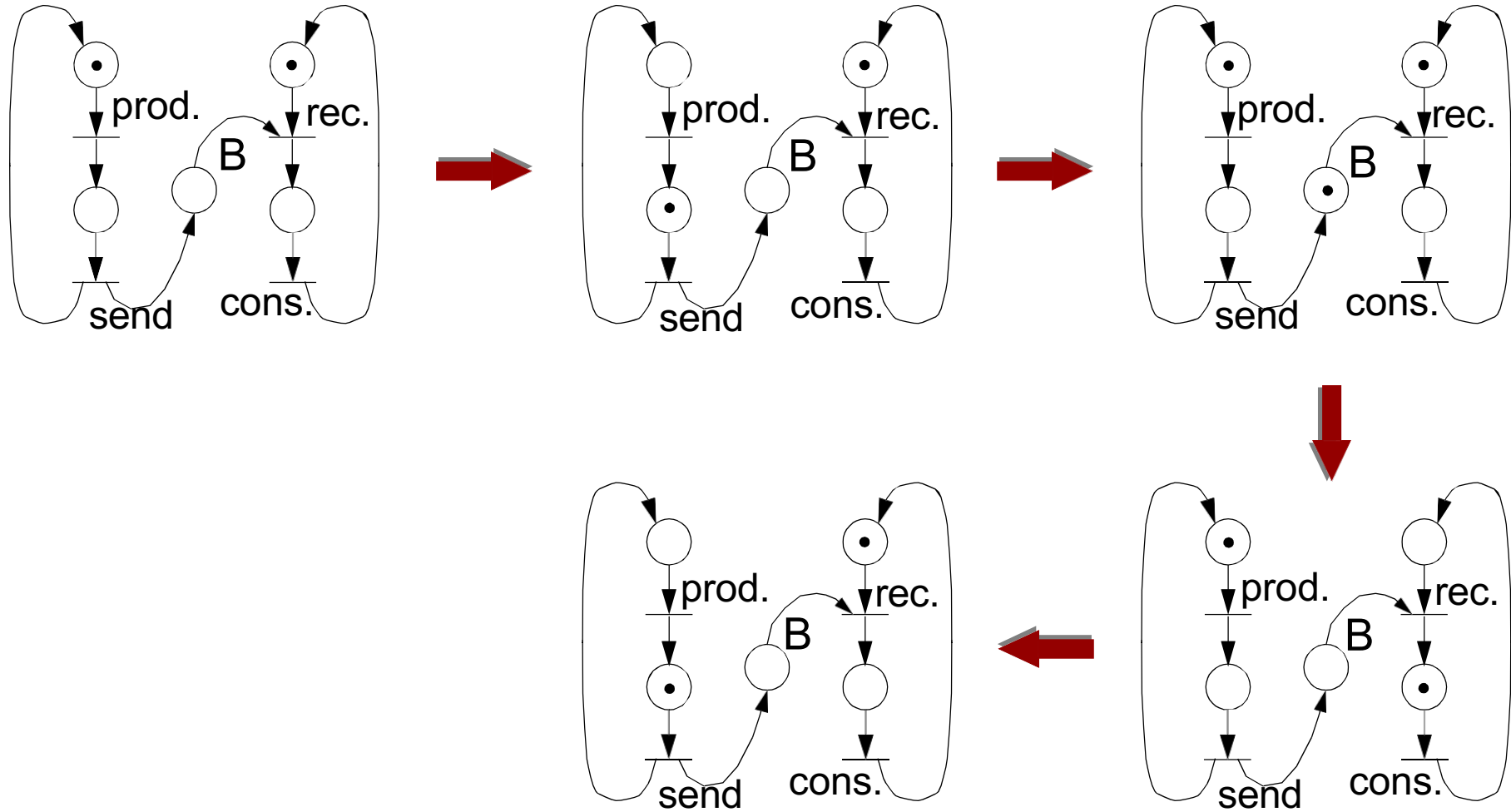
# Petri Net Example



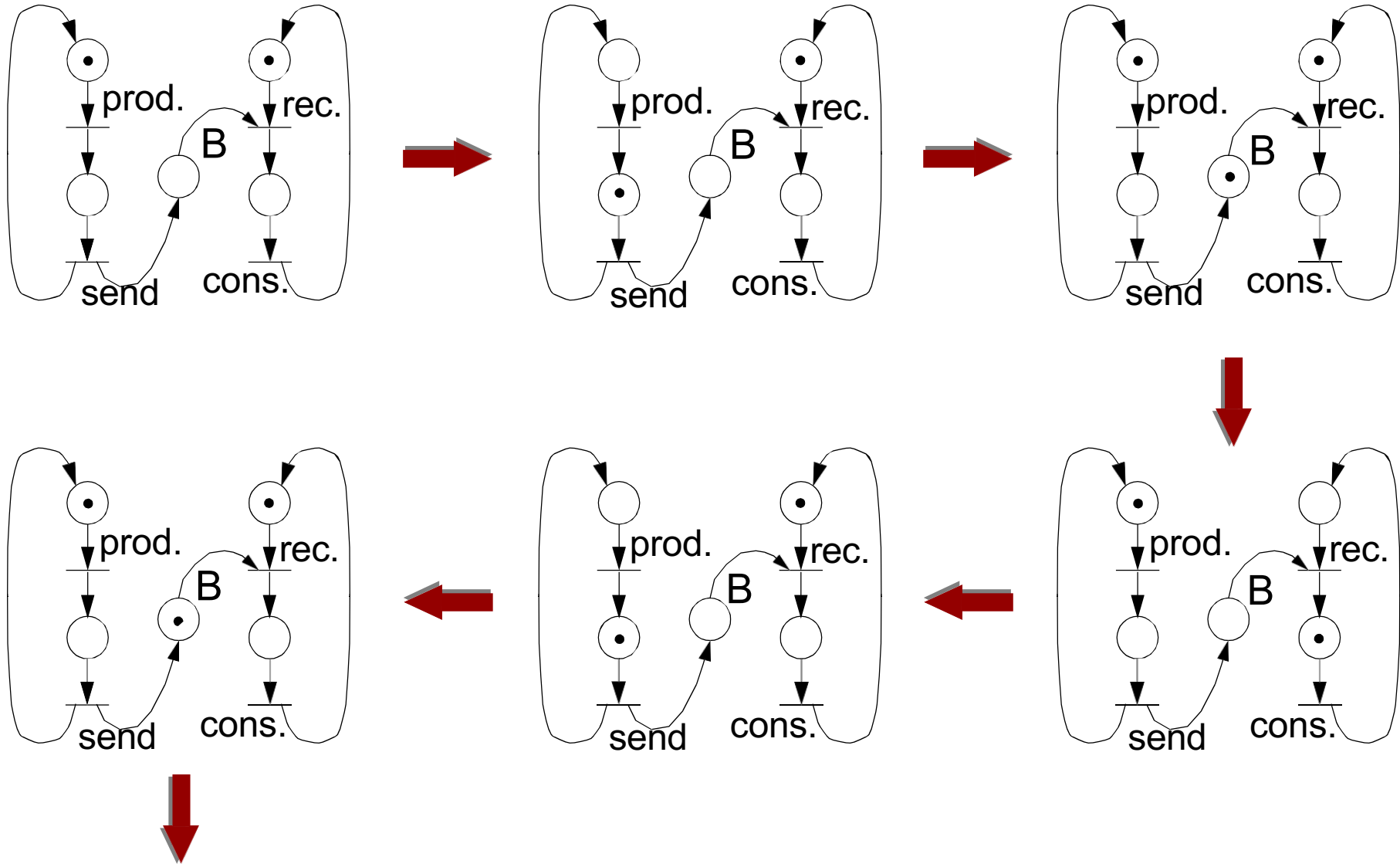
# Petri Net Example



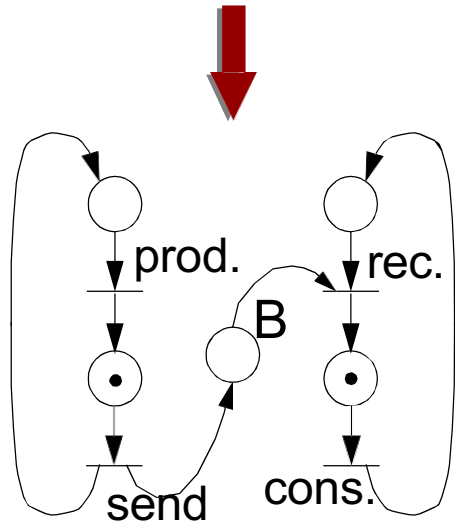
# Petri Net Example



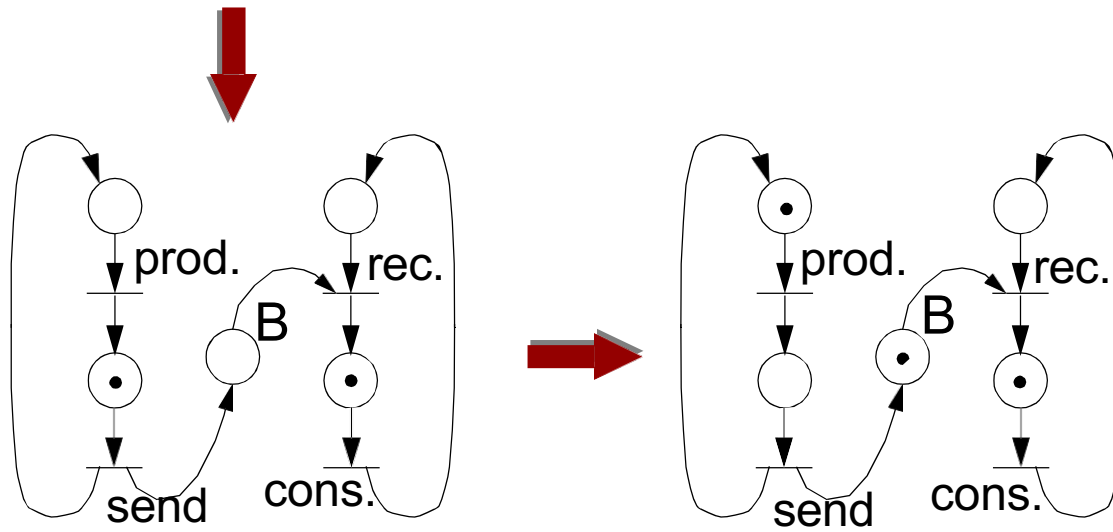
# Petri Net Example



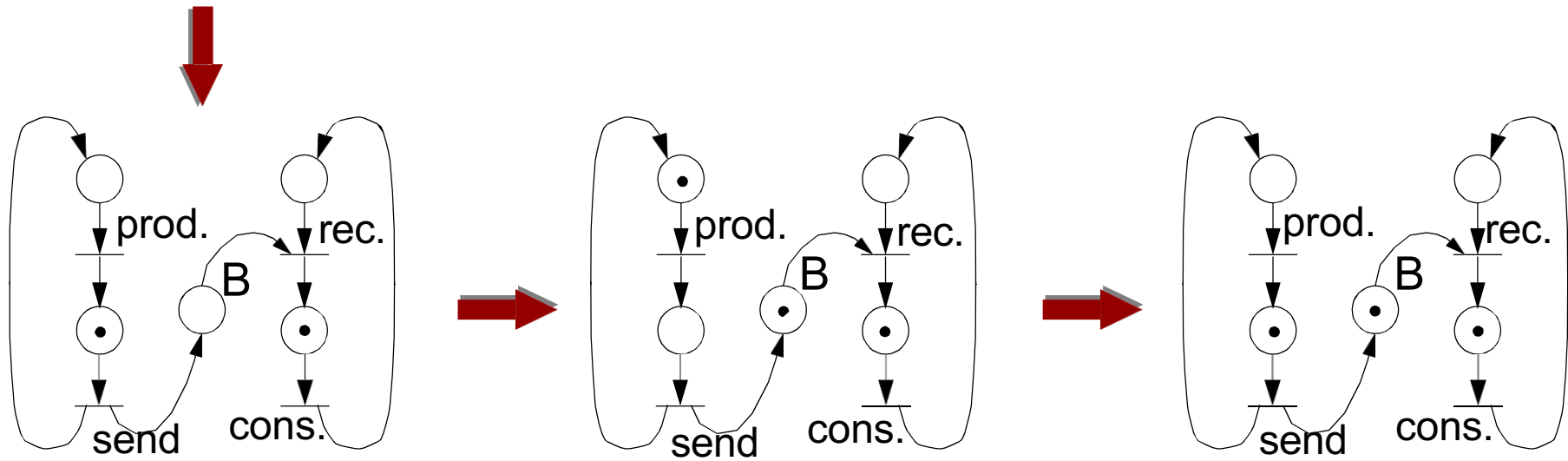
# Petri Net Example



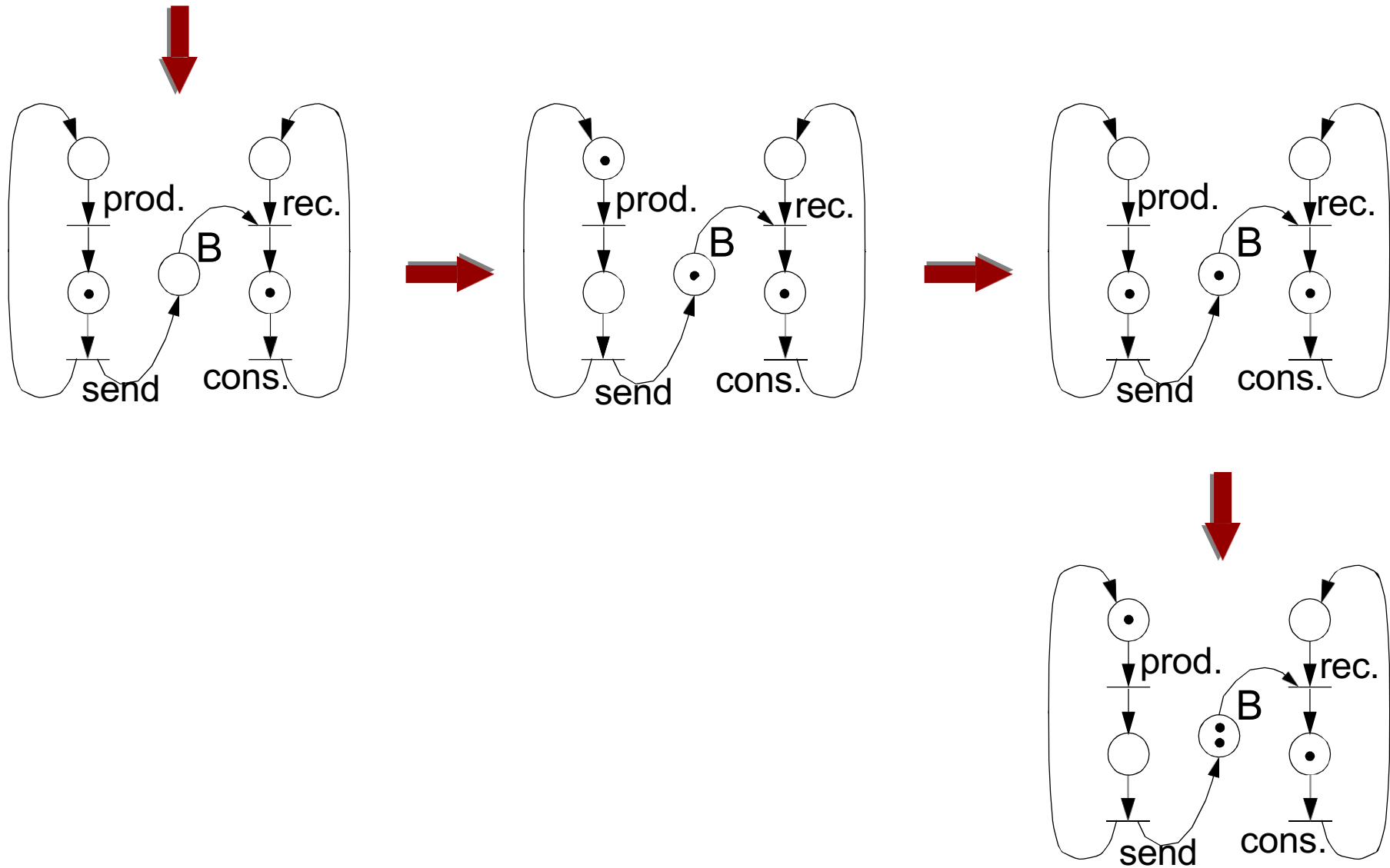
# Petri Net Example



# Petri Net Example

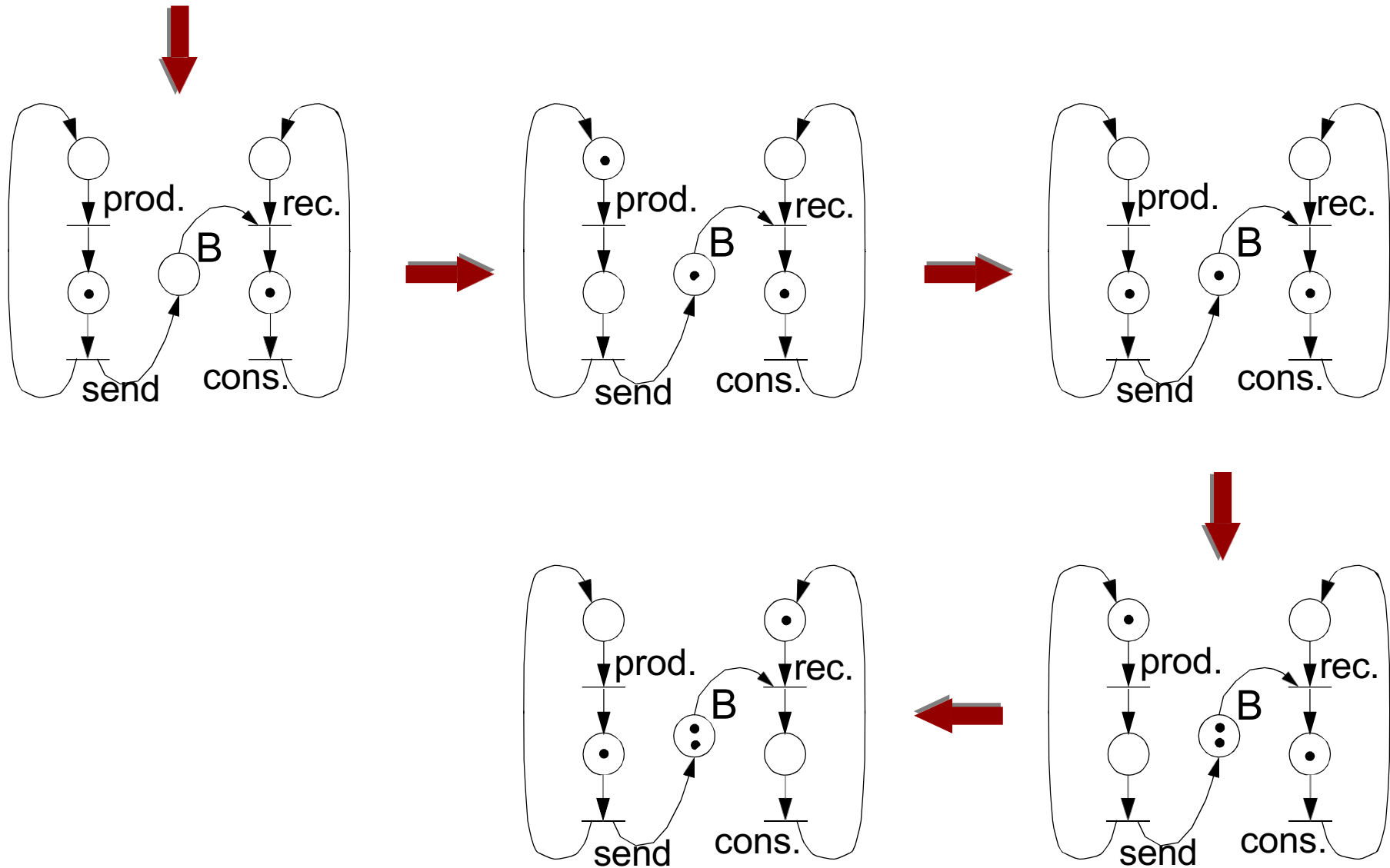


# Petri Net Example

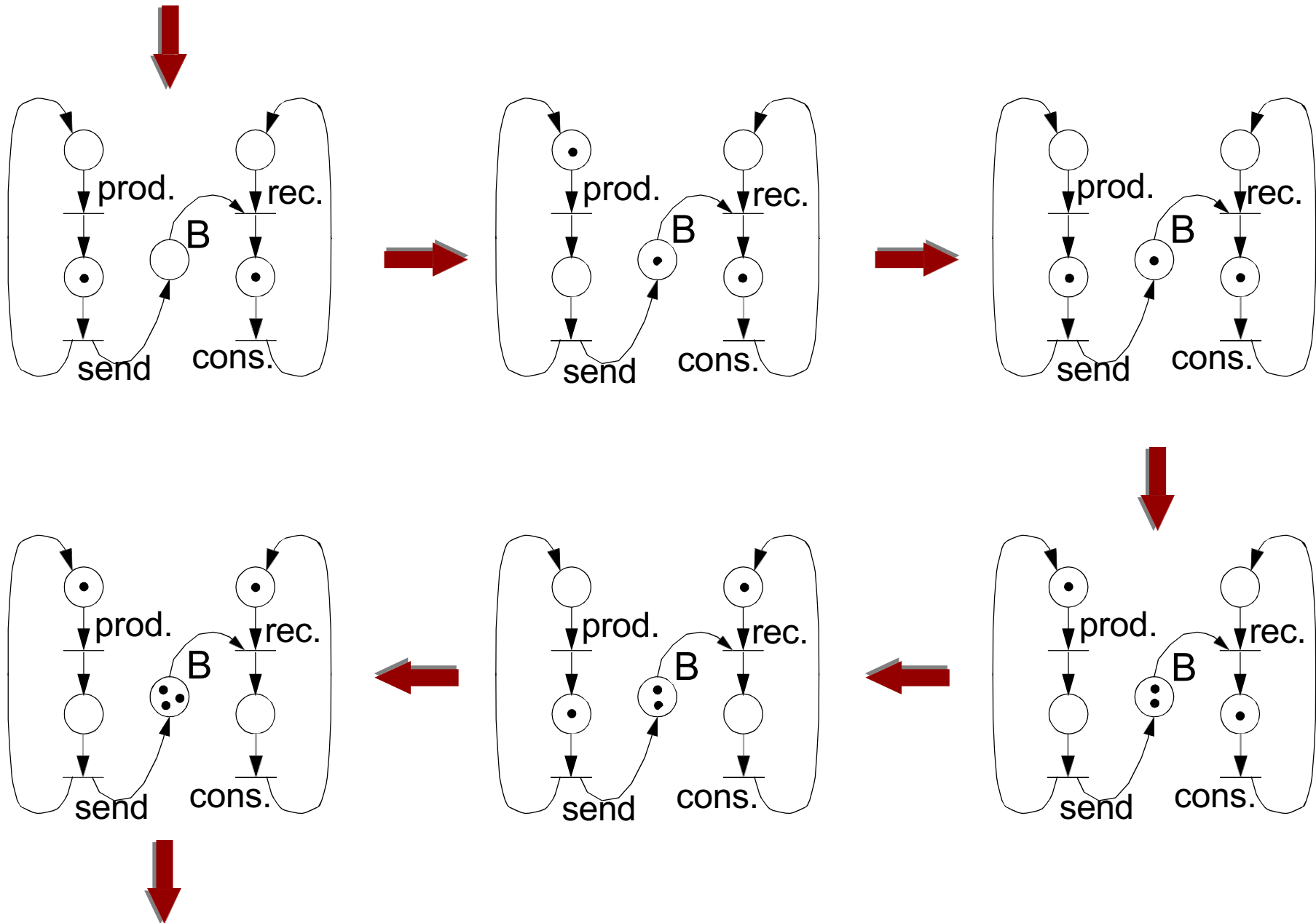




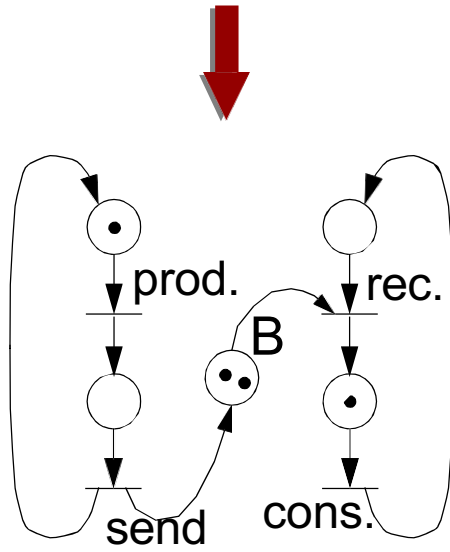
# Petri Net Example



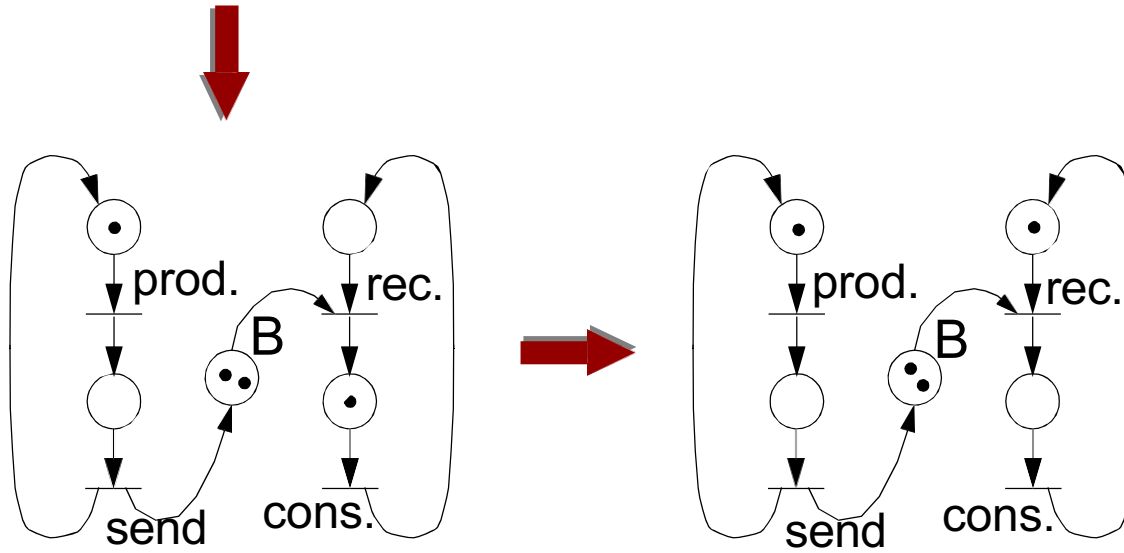
# Petri Net Example



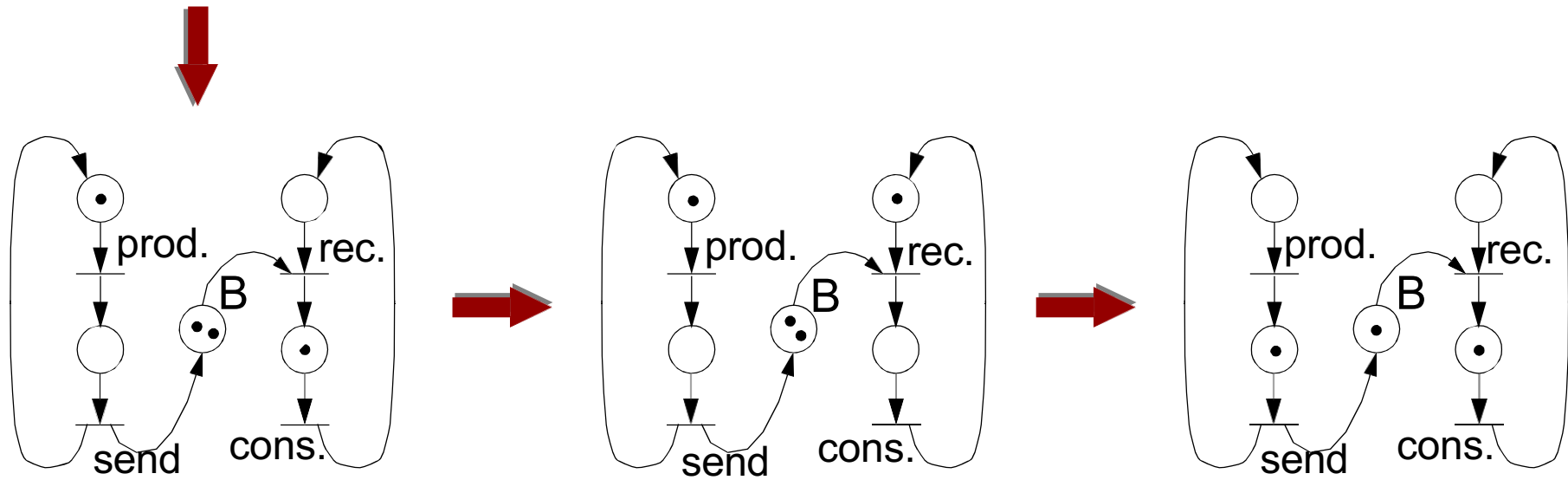
# Petri Net Example



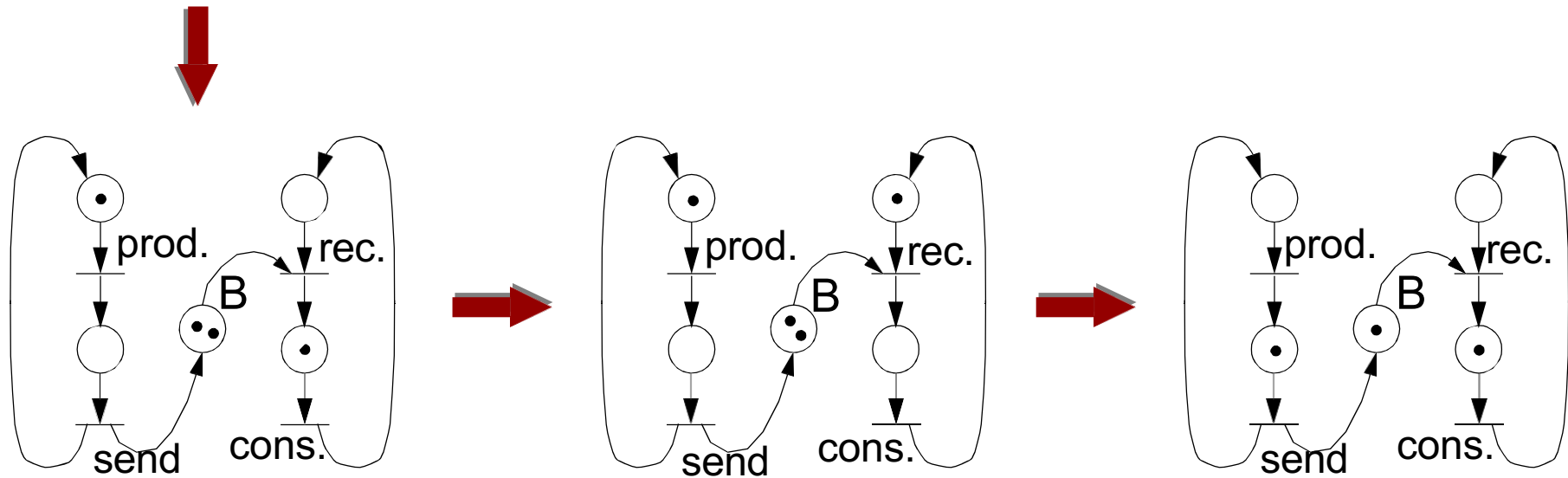
# Petri Net Example



# Petri Net Example



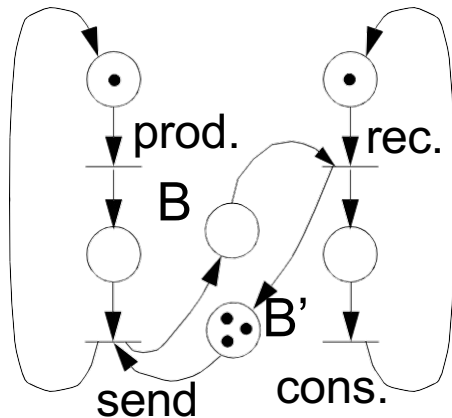
# Petri Net Example



- Notice that the buffer is considered to be infinite (tokens accumulate in *B*).

# Petri Net Example

Here we have the same model as on the previous slides, but with limited buffer.  
The buffer size is three (number of initial tokens in  $B'$ )

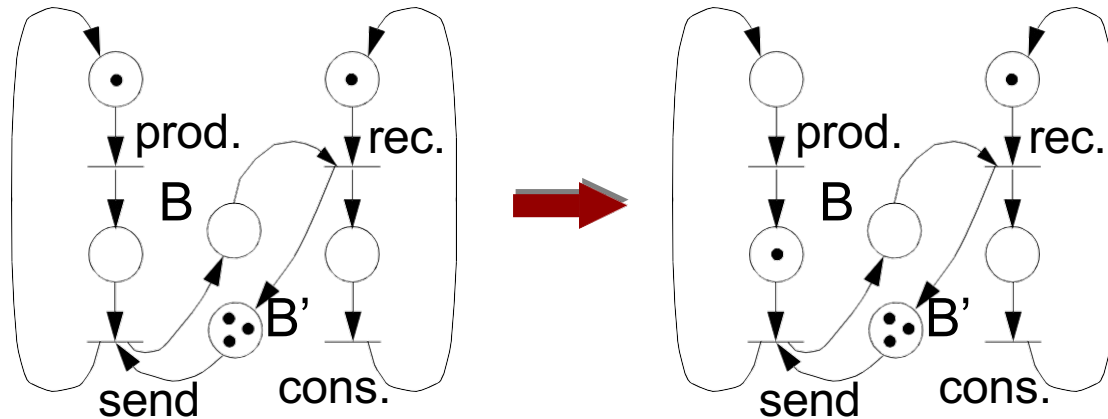


- Nr. of tokens in  $B'$ : how many free slots are available in the buffer;
- Nr. of tokens in  $B$ : how many messages (tokens) are in the buffer.

Total number of tokens in  $B$  and  $B'$  is constant (= 3).

# Petri Net Example

Here we have the same model as on the previous slides, but with limited buffer. The buffer size is three (number of initial tokens in  $B'$ )



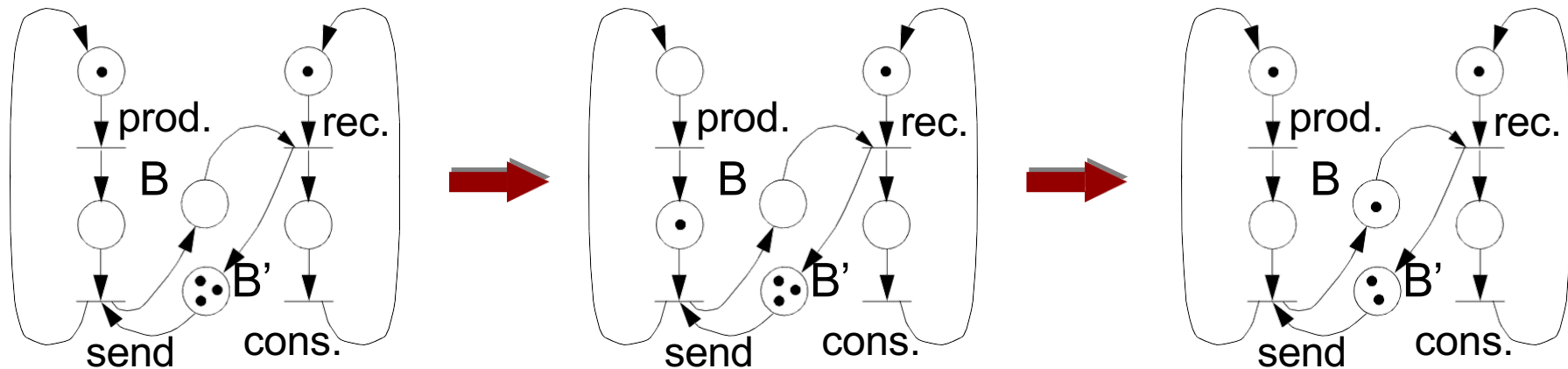
- Nr. of tokens in  $B'$ : how many free slots are available in the buffer;
- Nr. of tokens in  $B$ : how many messages (tokens) are in the buffer.

Total number of tokens in  $B$  and  $B'$  is constant (= 3).



# Petri Net Example

Here we have the same model as on the previous slides, but with limited buffer. The buffer size is three (number of initial tokens in  $B'$ )

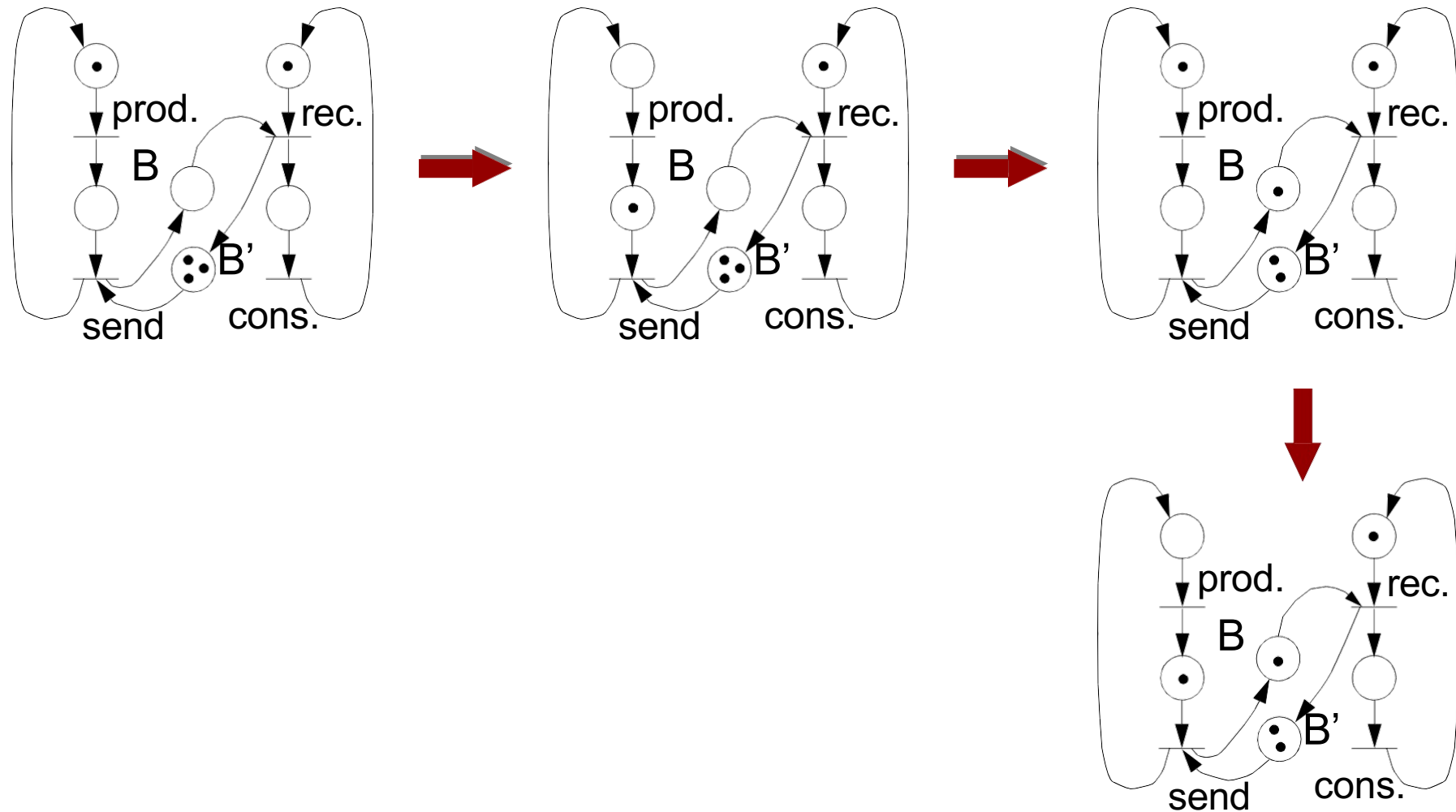


- Nr. of tokens in  $B'$ : how many free slots are available in the buffer;
- Nr. of tokens in  $B$ : how many messages (tokens) are in the buffer.

Total number of tokens in  $B$  and  $B'$  is constant (= 3).

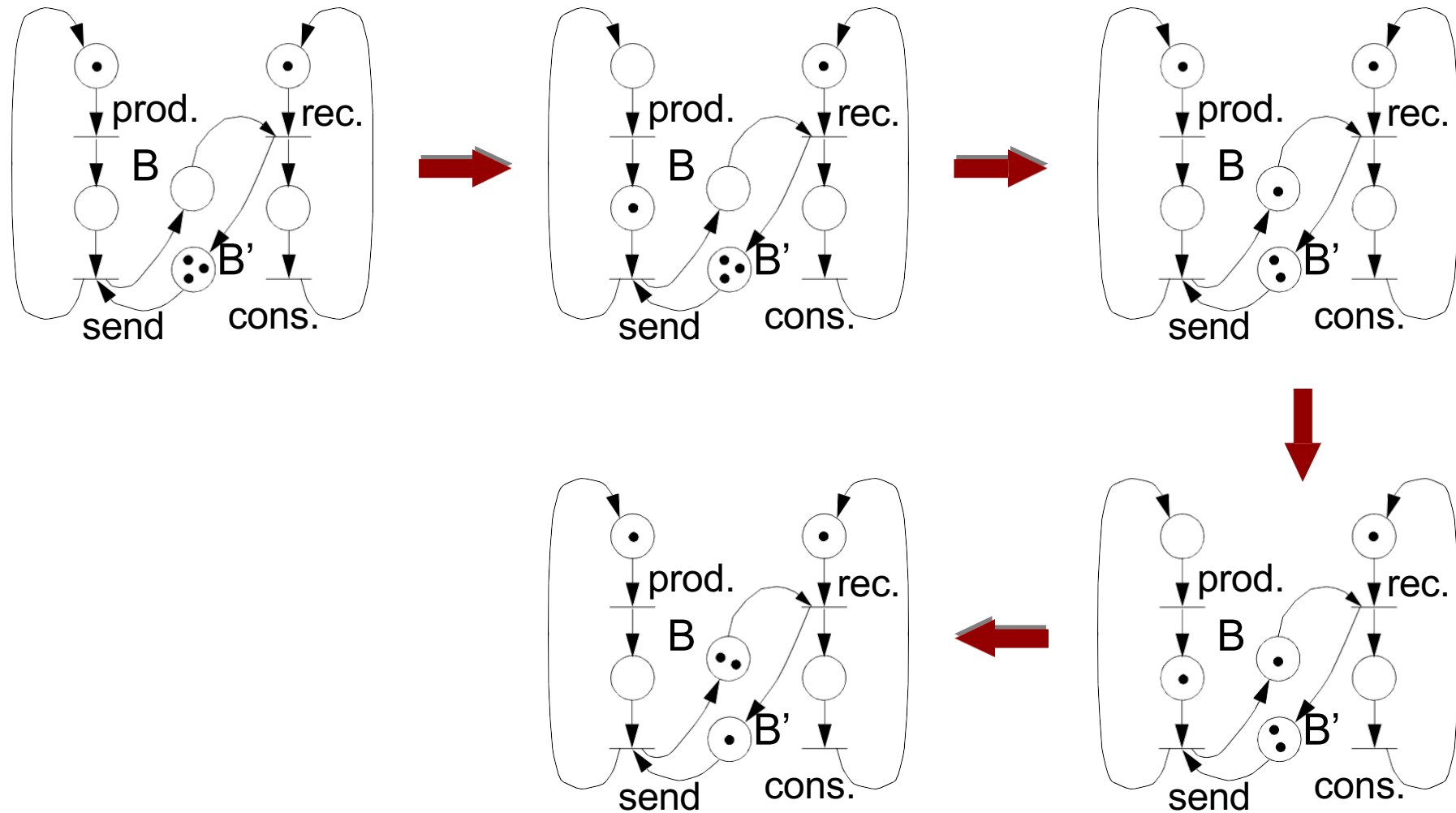
# Petri Net Example

Here we have the same model as on the previous slides, but with limited buffer. The buffer size is three (number of initial tokens in  $B'$ )



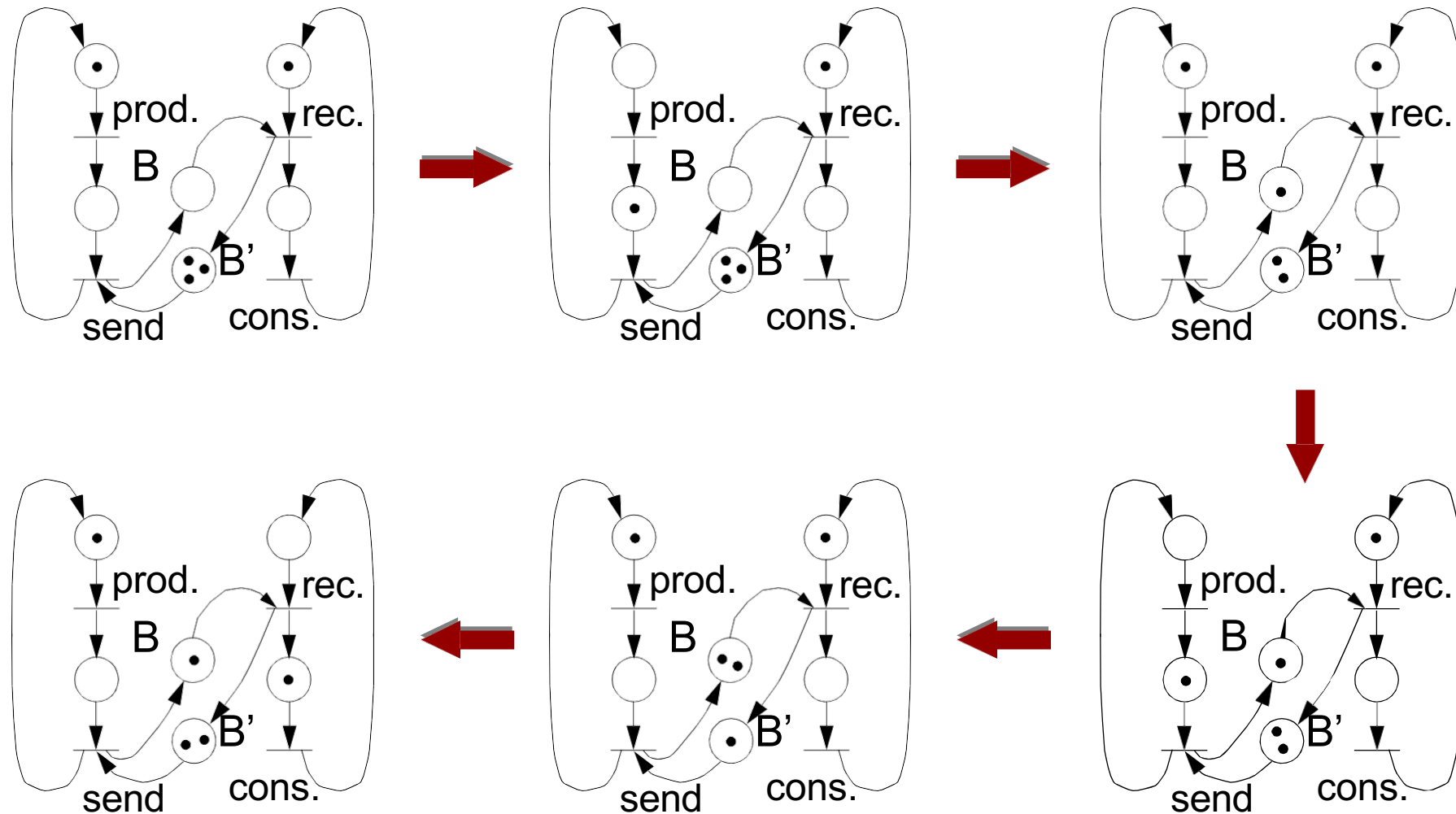
# Petri Net Example

Here we have the same model as on the previous slides, but with limited buffer. The buffer size is three (number of initial tokens in  $B'$ )



# Petri Net Example

Here we have the same model as on the previous slides, but with limited buffer. The buffer size is three (number of initial tokens in  $B'$ )



# Some Features and Applications of Petri Nets

- **Intuitive.**

**Easy to express concurrency, synchronisation, nondeterminism.**

***Nondeterminism is an important difference between Petri nets and dataflow!***

- **As an uninterpreted model, Petri Nets can be used for several, very different classes of problems.**

- ***Uninterpreted model:* nothing has to be specified related to the particular activities associated to the transitions.**

# Some Features and Applications of Petri Nets

- **Petri Nets have been intensively used for modeling and analysis of industrial production systems, information systems, but also**
  - **Computer architectures**
  - **Operating systems**
  - **Concurrent programs**
  - **Distributed systems**
  - **Hardware systems**

# Properties and Analysis of Petri Nets

- Several properties of the system can be analysed using Petri nets:
  - **Boundedness:** number of tokens in a place does not exceed a limit.  
If this limit is 1, the property is sometimes called *safeness*.
    - You can check that available resources are not exceeded.

# Properties and Analysis of Petri Nets

- Several properties of the system can be analysed using Petri nets:
  - **Boundedness:** number of tokens in a place does not exceed a limit.  
If this limit is 1, the property is sometimes called *safeness*.
    - You can check that available resources are not exceeded.
  - **Liveness:** A transition  $t$  is called live if for every possible marking there exists a chance for that transition to become enabled.  
The whole net is live, if all its transitions are live.
    - Important in order to check that the system is not deadlocked.



# Properties and Analysis of Petri Nets

- Several properties of the system can be analysed using Petri nets:
  - **Boundedness:** number of tokens in a place does not exceed a limit.  
If this limit is 1, the property is sometimes called *safeness*.
    - You can check that available resources are not exceeded.
  - **Liveness:** A transition  $t$  is called live if for every possible marking there exists a chance for that transition to become enabled.  
The whole net is live, if all its transitions are live.
    - Important in order to check that the system is not deadlocked.
  - **Reachability:** given a current marking  $M$  and another marking  $M'$ , does there exist a sequence of transitions by which  $M'$  can be obtained?
    - You can check that a certain desired state (marking) is reached.
    - You can check that a certain undesired state is never reached.

# Properties and Analysis of Petri Nets

Mathematical tools are available for analysis of Petri Nets.



The properties discussed above can be formally verified.

- Petri nets (like dataflow systems) are *asynchronous concurrent*.
  - Events can happen at any time.
  - There exists a partial order of events.

# Extended Petri Net Models

Basic Petri Net models have a limited expressive power.

# Extended Petri Net Models

Basic Petri Net models have a limited expressive power.

- **Timed Petri Nets**

- Transitions have associated times (time intervals)
- Tokens are carrying time stamps.

With timed Petri nets we can model the timing aspects

# Extended Petri Net Models

Basic Petri Net models have a limited expressive power.

## ■ Timed Petri Nets

- Transitions have associated times (time intervals)
- Tokens are carrying time stamps.

With timed Petri nets we can model the timing aspects

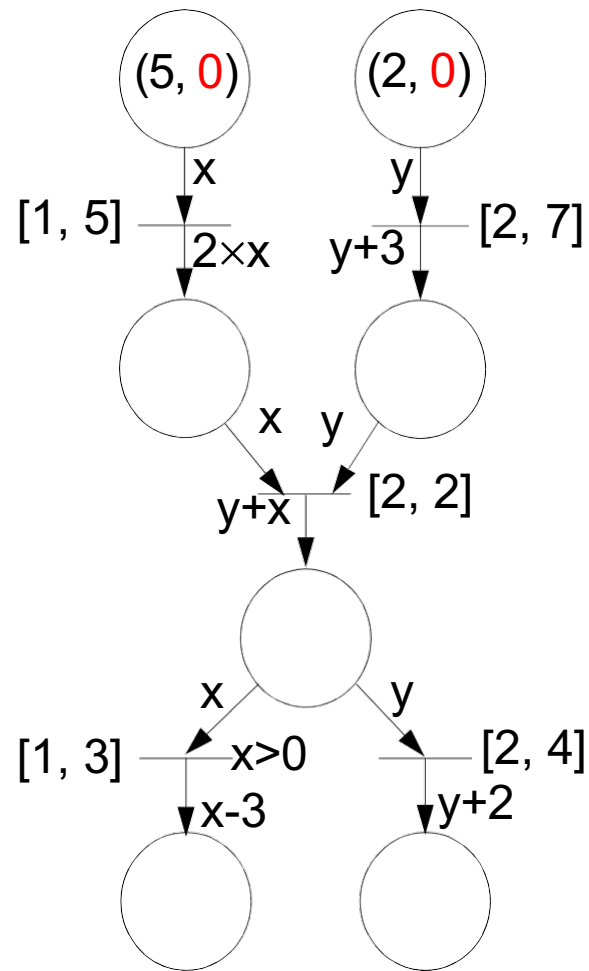
## ■ Coloured Petri Nets

- Tokens have associated values
- Transitions have associated functions

Coloured Petri Nets are similar to dataflow models (but also capture nondeterminism!).

# Extended Petri Net Models

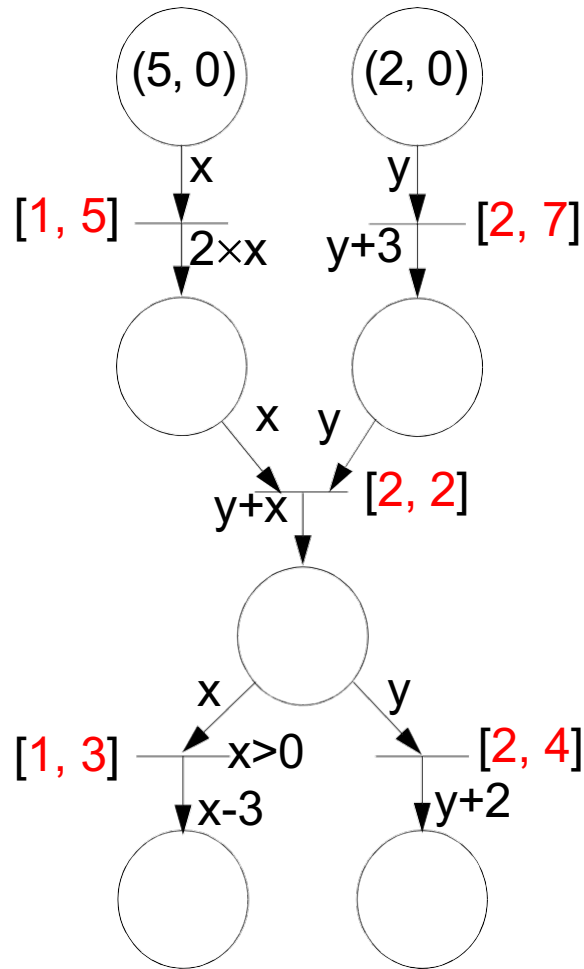
## Coloured and Timed Petri net



- Tokens carry Time stamps

# Extended Petri Net Models

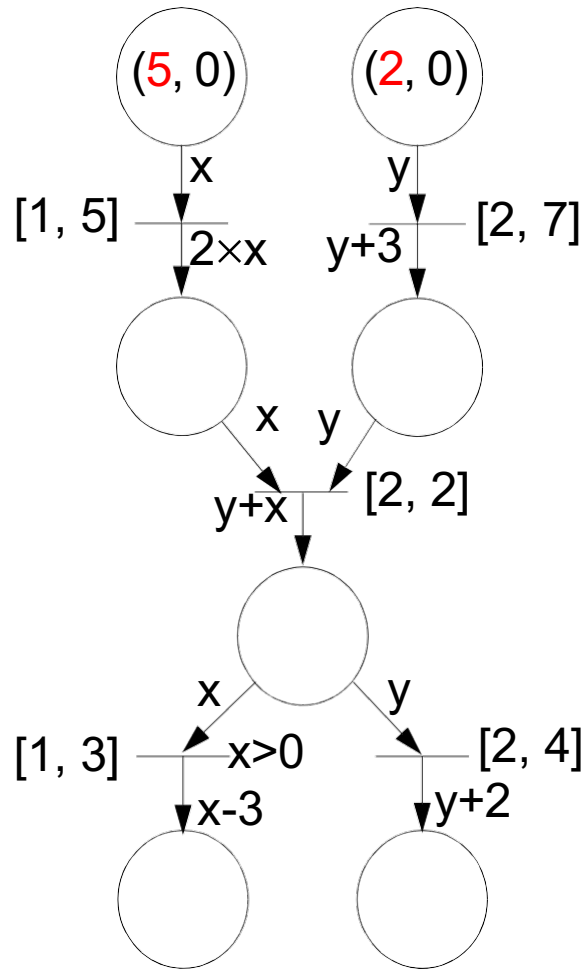
## Coloured and Timed Petri net



- Tokens carry Time stamps
- **Transitions have associated time (interval)**

# Extended Petri Net Models

## Coloured and Timed Petri net

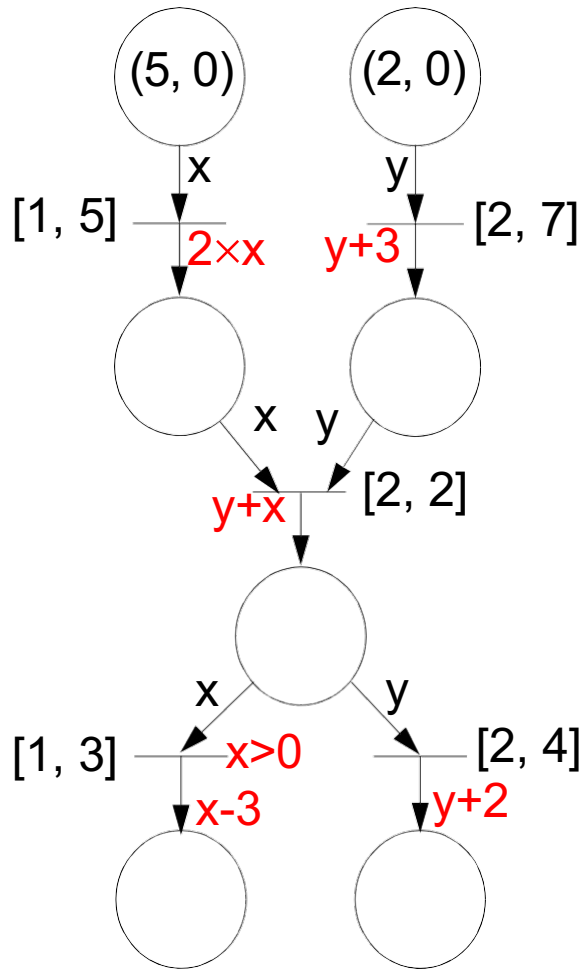


- Tokens carry Time stamps
- Transitions have associated time (interval)
- Tokens have associated values



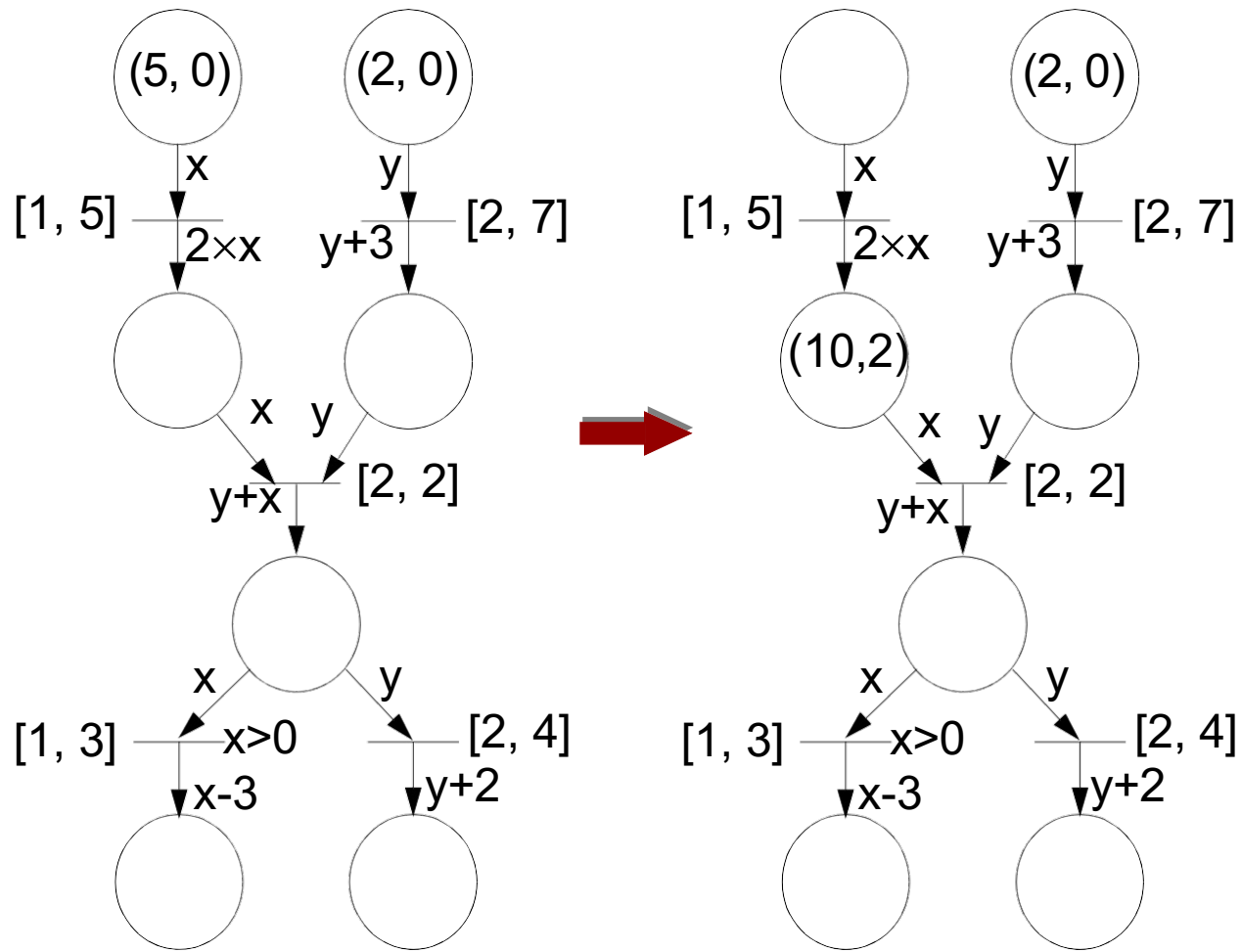
# Extended Petri Net Models

## Coloured and Timed Petri net

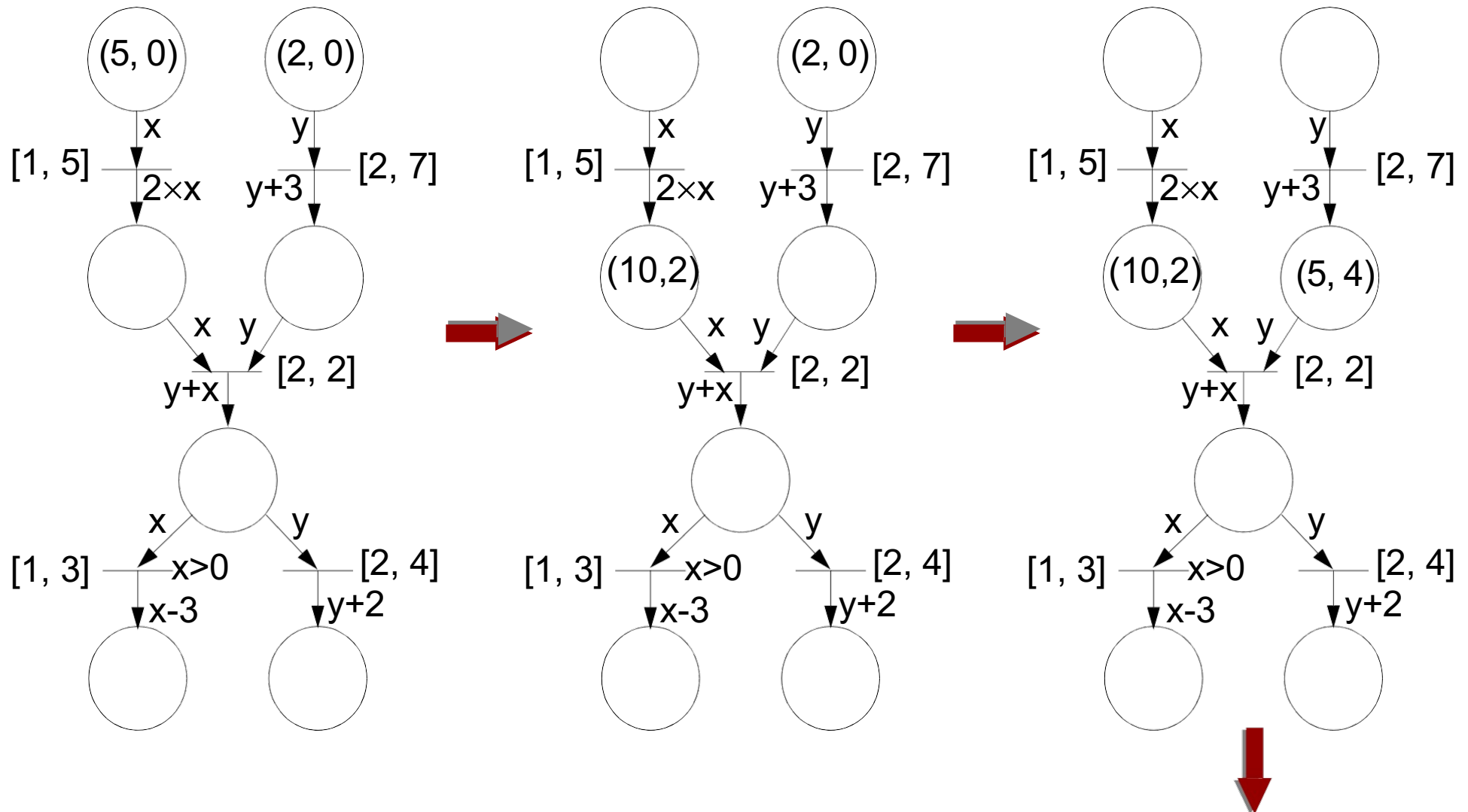


- Tokens carry Time stamps
- Transitions have associated time (interval)
- Tokens have associated values
- Transitions have associated functions and guards

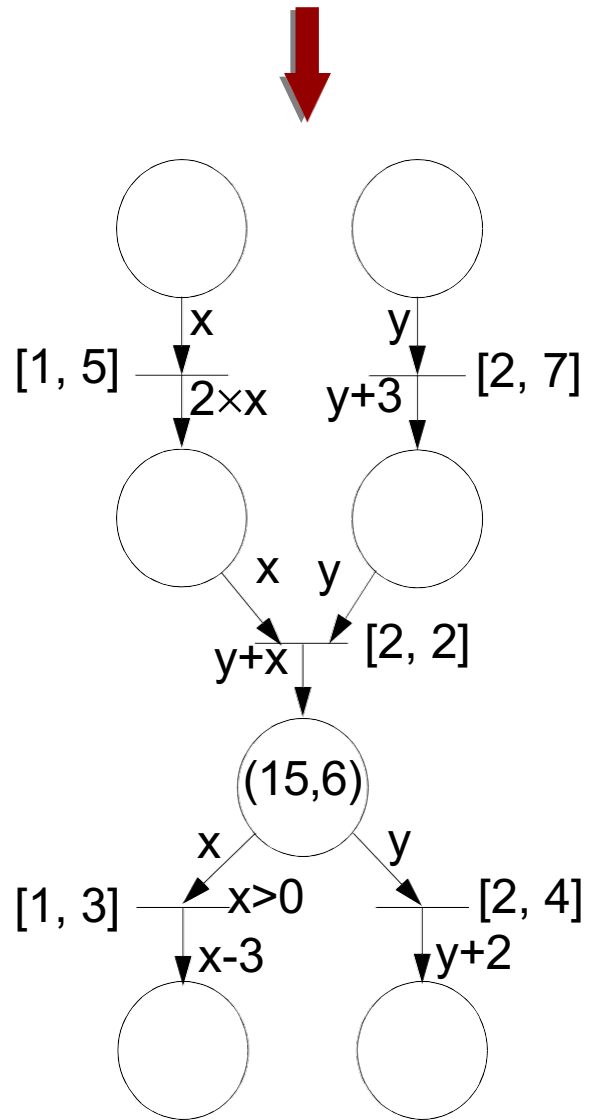
# Extended Petri Net Models



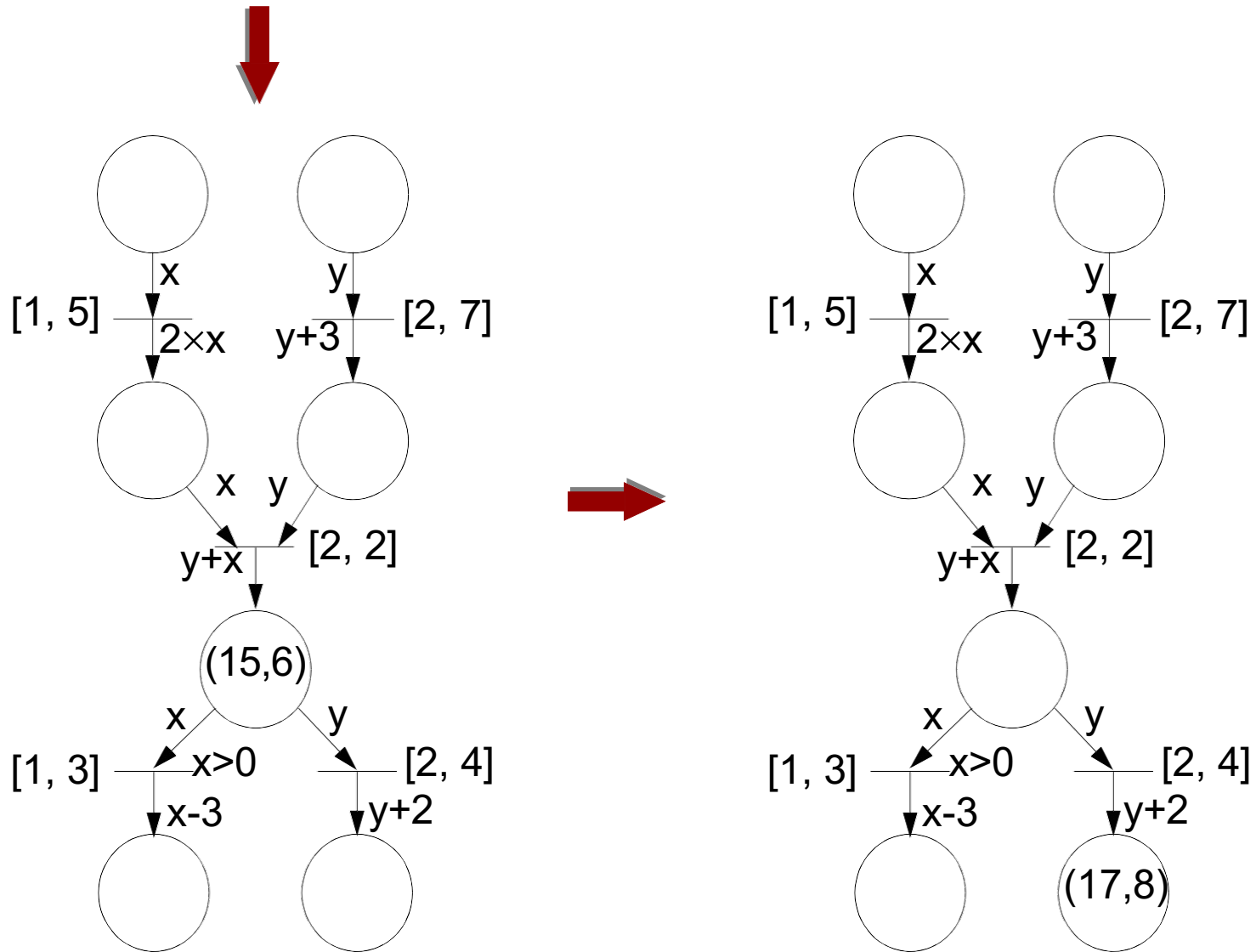
# Extended Petri Net Models



# Extended Petri Net Models



# Extended Petri Net Models



# Extended Petri Net Models

- Extended Petri Nets have a larger expressive power than classical Petri Nets.



Analysis is more complex; the formal analysis of properties can take very large amounts of time (memory).

- Simulation of the Petri Net is very often used in order to verify the system and to estimate performance